

# Improving the Data Delivery Latency in Sensor Networks with Controlled Mobility

Seminar: Ad-hoc networks  
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# Outline

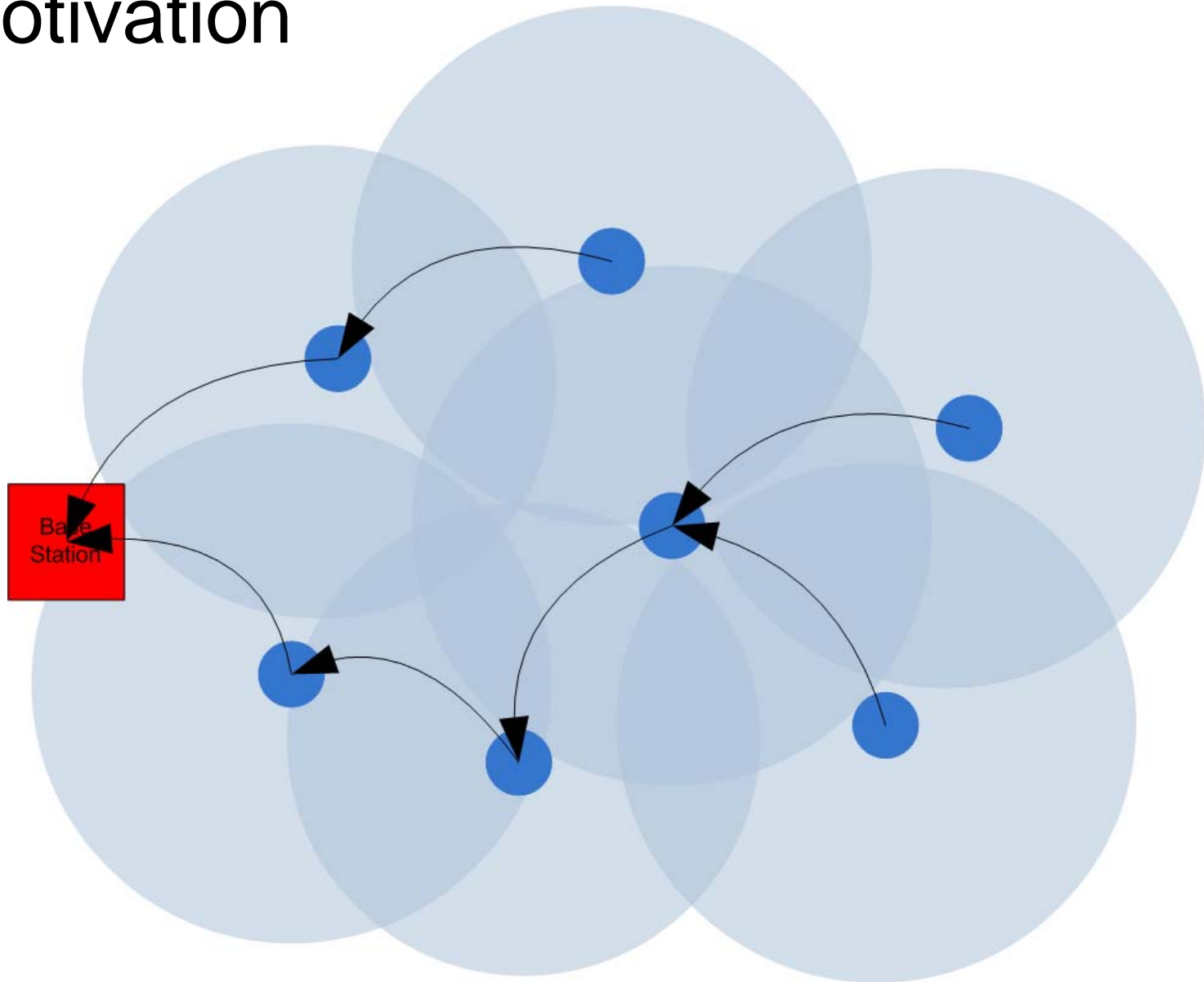
1. Motivation
2. Path Selection Problem
3. Approximation Algorithm
4. Application Scenario
5. Experiments

# Motivation

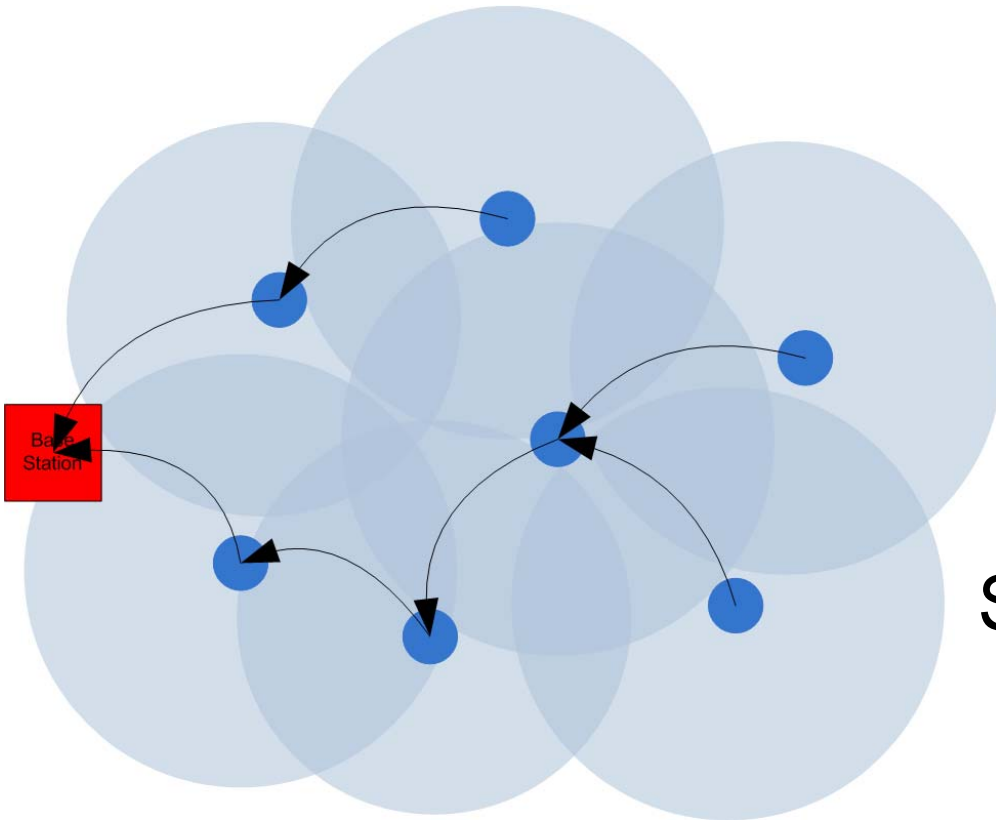


Base  
Station

# Motivation



# Motivation – Energy Problem

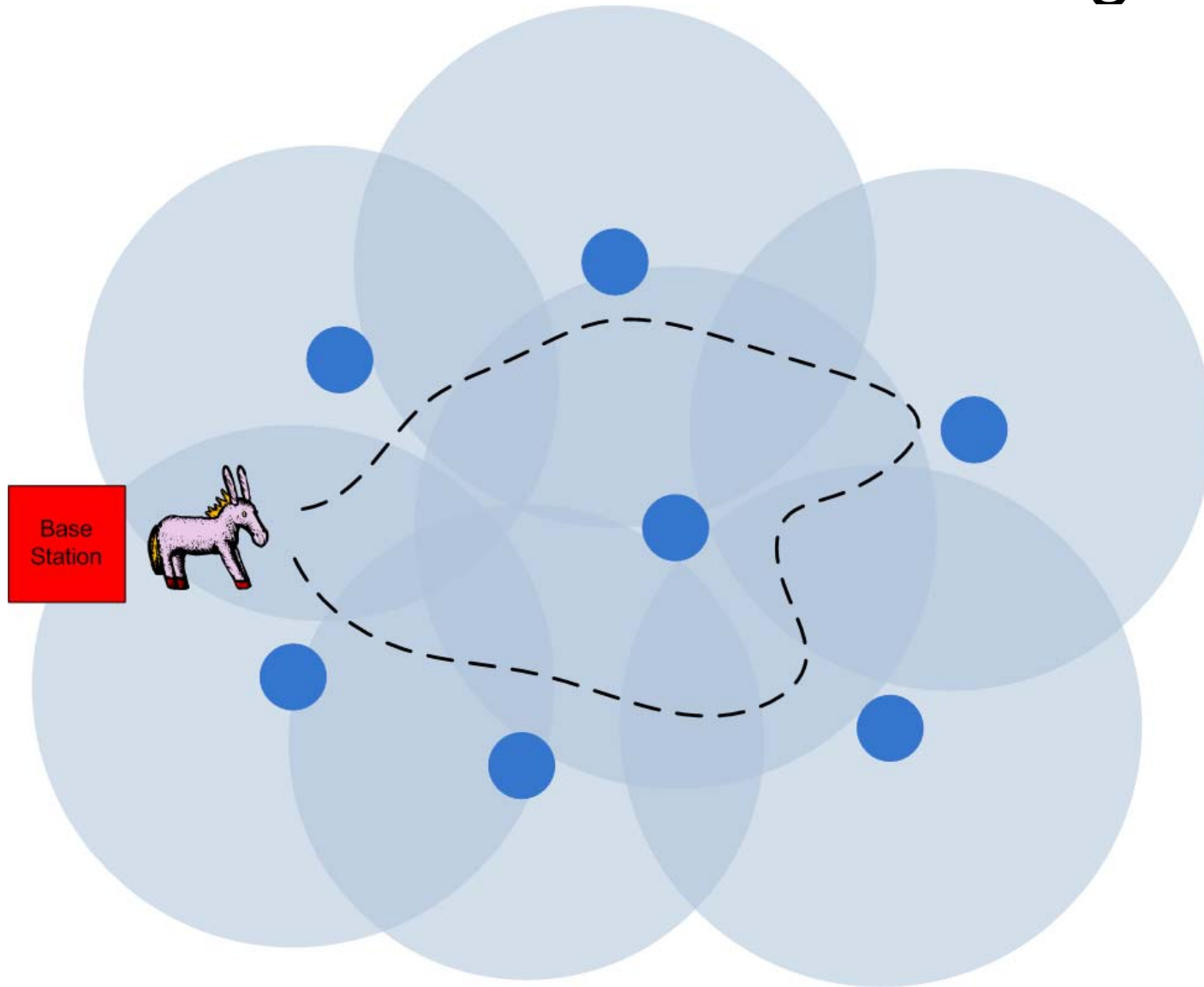


- Complex algorithms
- Continual data forwarding
- Long communication ranges

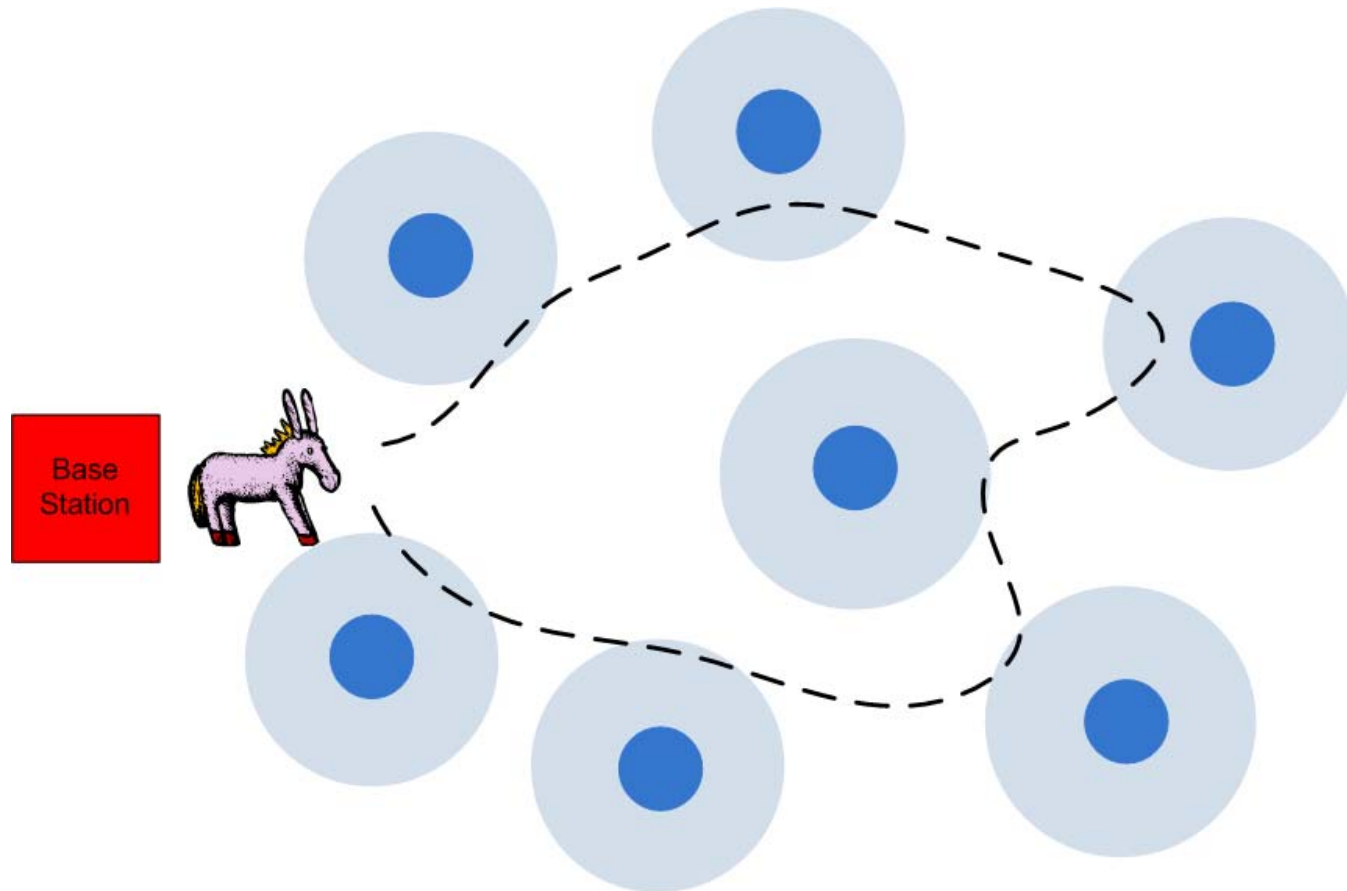
## Solution:

Mobile data collector  
(Data Mule)

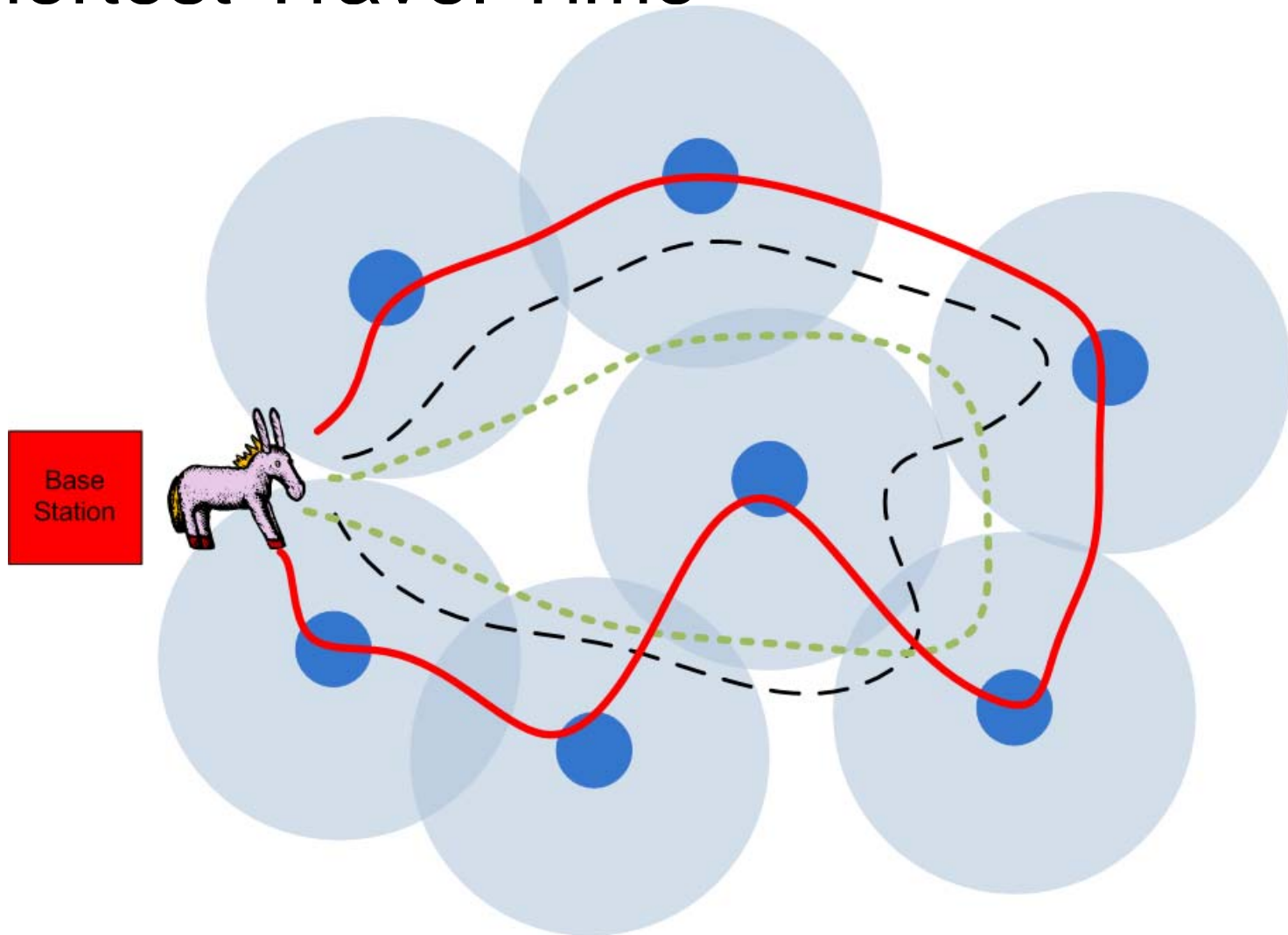
# Motivation – No Data Forwarding



# Motivation – Short Communication Ranges

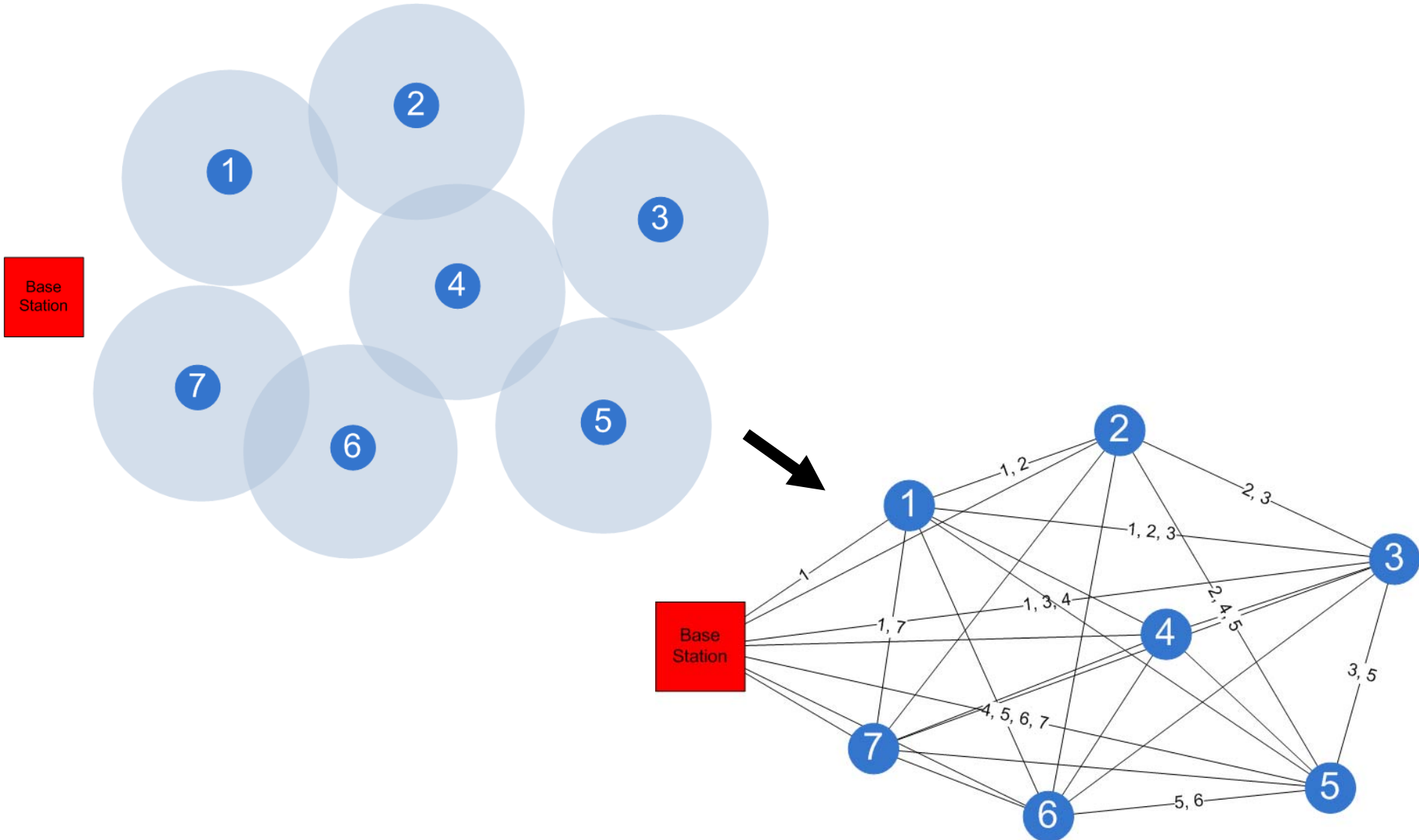


# Path Selection Problem – Shortest Travel Time

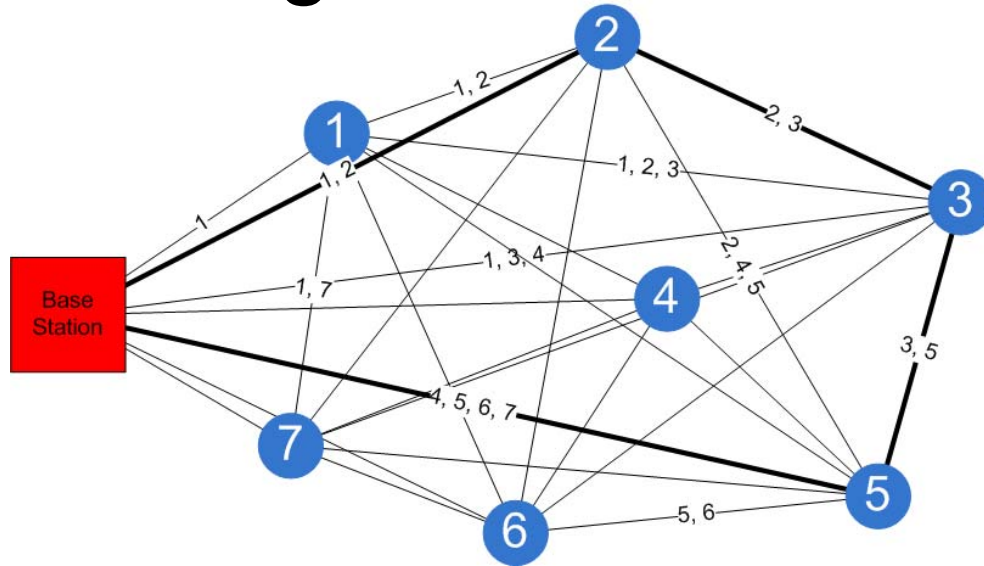




# Simplifying the Problem



# Label-Covering Tour



## Definition:

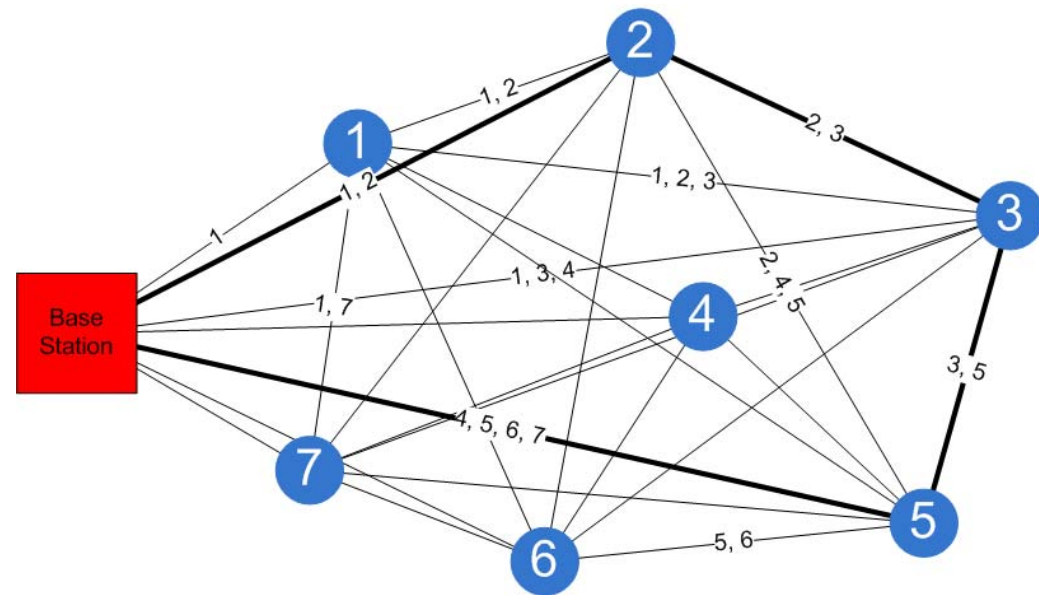
- vertices are points in  $R^2$  (static nodes)
- edges are possible paths
- labels on edges indicate possible communication of nodes
- cost for each edge to indicate travel time (not shown)

## Problem:

Find a set of edges that covers all labels and minimizes the costs.

# Choice of Cost Metric

- Number of edges
- Euclidean distance
- Uncovered distance



# Choice of Cost Metric

- Creation of different node deployments
- Simulation of a random walk
- Calculate correlation to the total travel time

Radius ( $d$ )	150				500			
Comm. range ( $r$ )	10		100		10		100	
Exec. time ( $e$ )	2	20	2	20	2	20	2	20
Num. edge	0.992	0.987	0.982	0.850	0.984	0.982	0.988	0.988
Euclidean dist.	0.997	0.996	0.990	0.835	0.999	0.999	0.999	0.999
Uncovered dist.	0.992	0.993	—	—	0.999	0.999	0.935	0.935

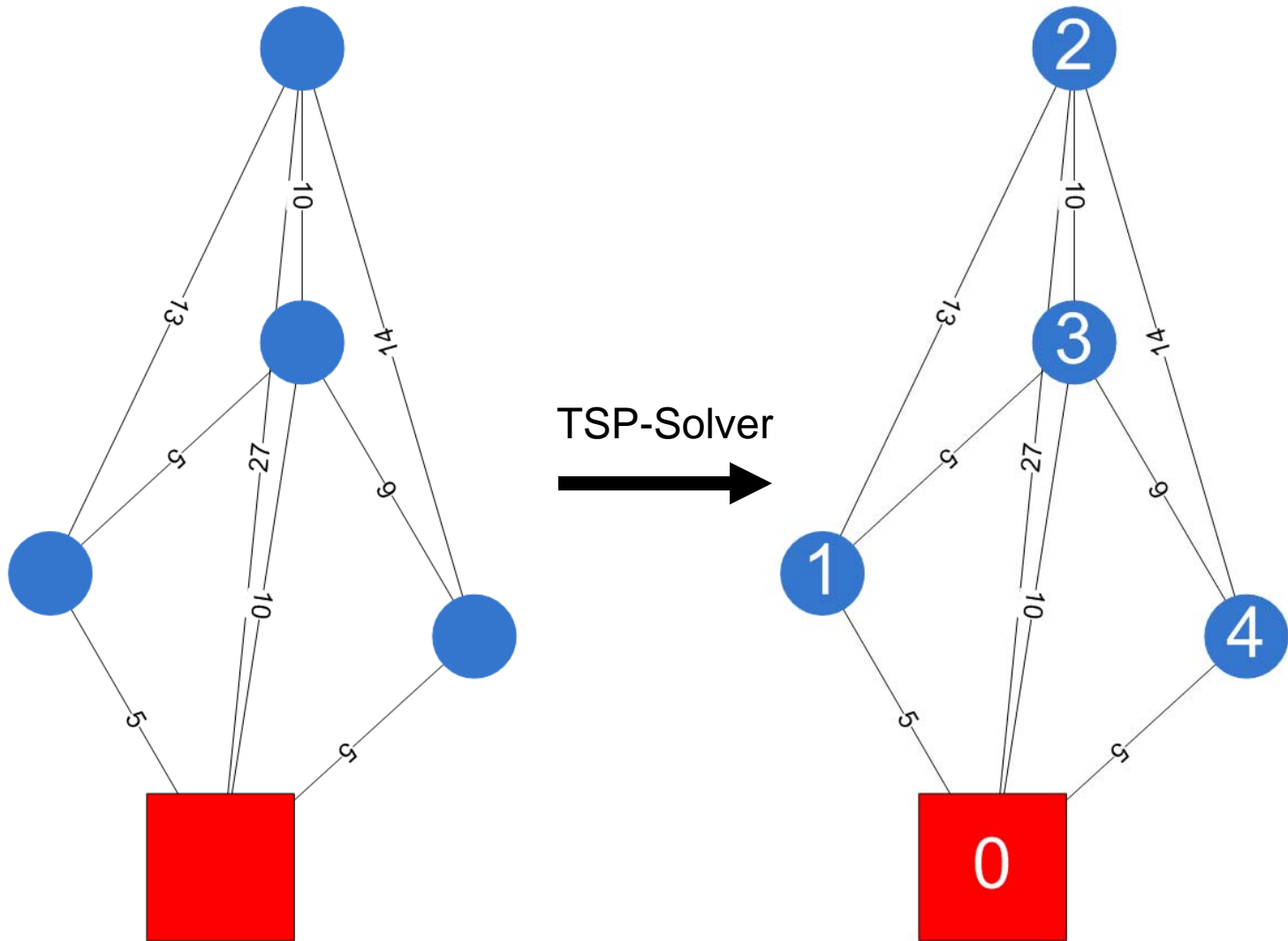


# Label-Covering Tour is NP-Hard

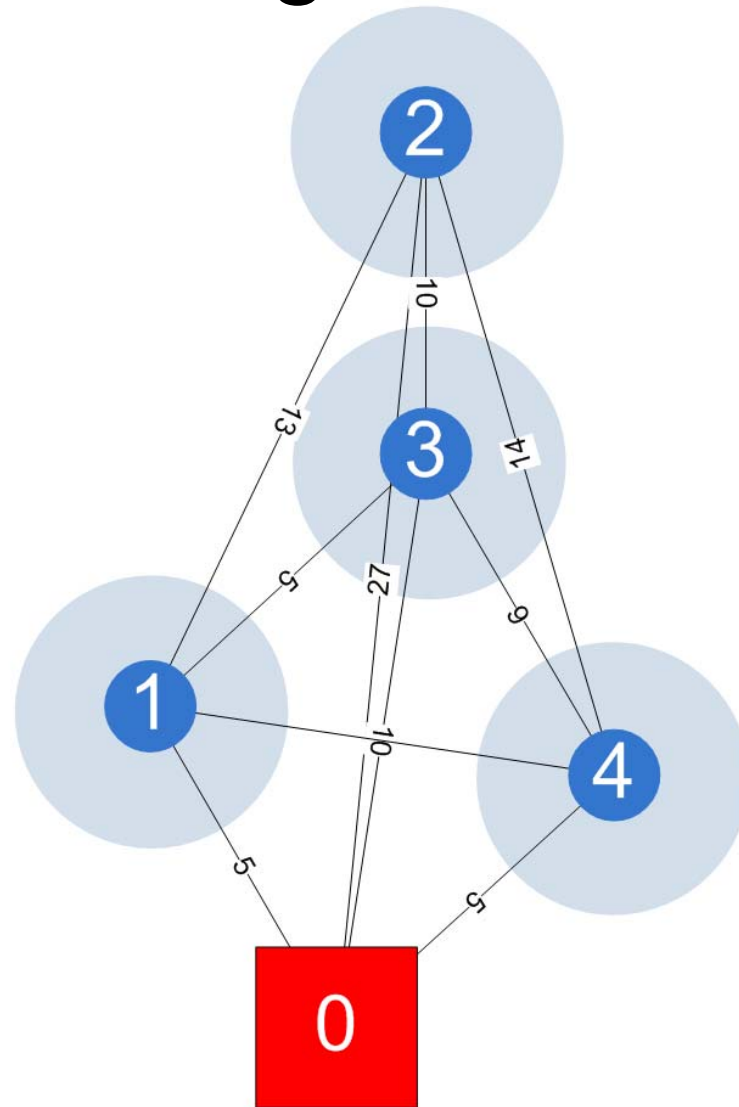
Choose a small communication range

- Max two labels on each edge
- Label-Covering Tour must visit all nodes
- Problem is equivalent to TSP

# Approximation Algorithm



# Approximation Algorithm





# Approximation Algorithm - Analysis

- Costs of TSP depends on the algorithm
- Costs of suggested Approx-Algorithm:  $O(n^3)$
- Finding Label-Covering Tour:  
 $C(\text{TSP}) + O(n^3)$



# Example Application

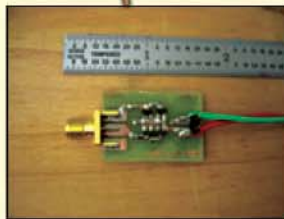
Peak displacement sensor



802.11 g and Zigbee wireless communication



Control GUI and data acquisition



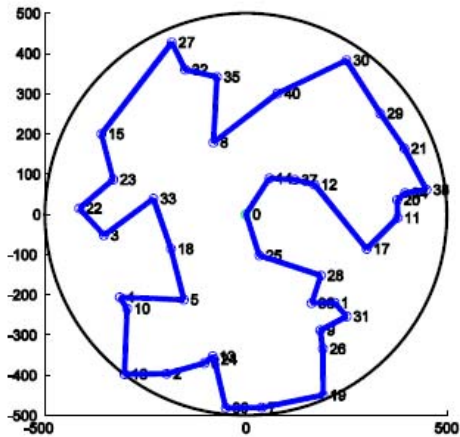
RF power delivery



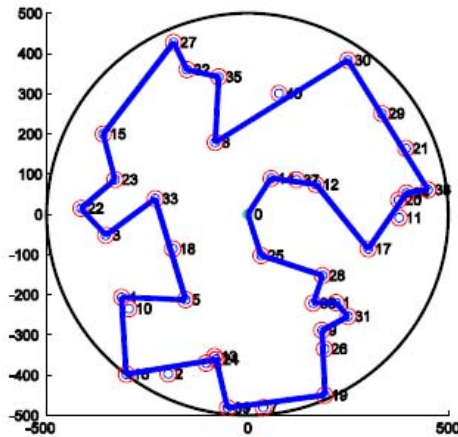
On-board computing and RF source



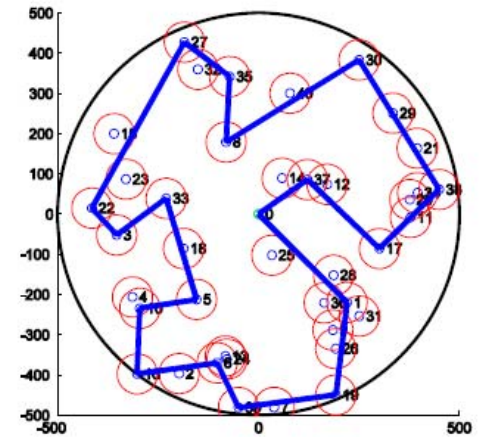
# Experiments



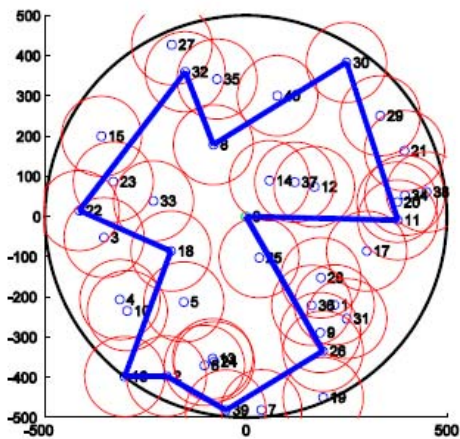
$r = 1$



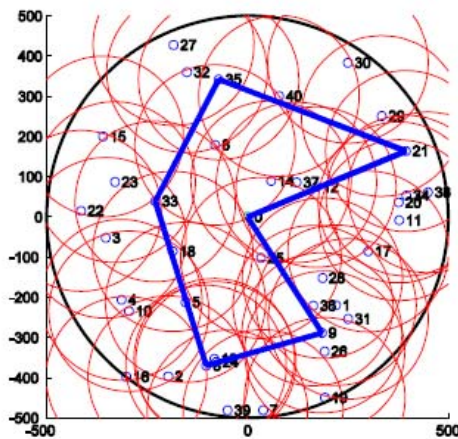
$r = 20$



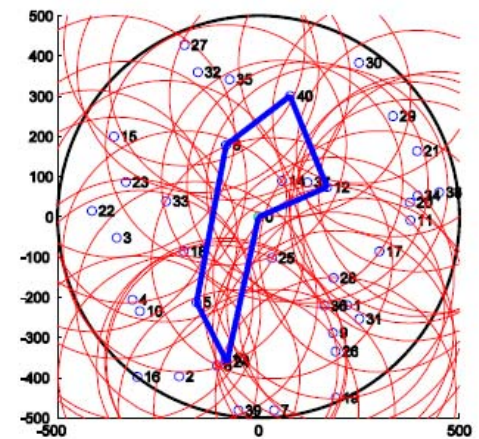
$r = 50$



$r = 100$



$r = 200$



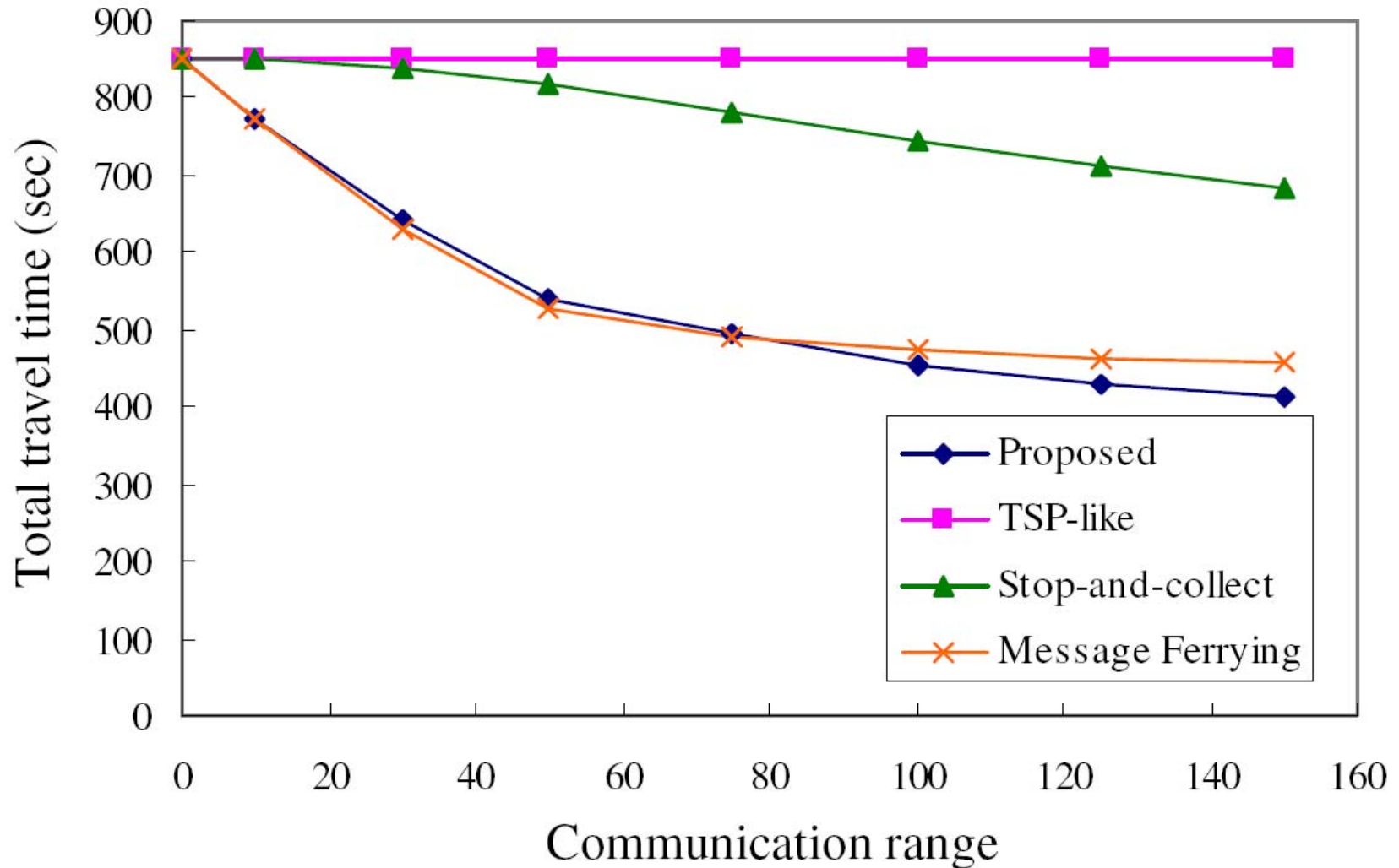
$r = 300$



# Comparison with Other Strategies

- TSP-like
- Stop-and-collect
- Message-Ferrying

# Comparison with Other Strategies





# Conclusion

- Limited energy is a problem in sensor networks
- Mobility for data collection prolongs network lifetime
- Path selection is essential for data mules
- Approximation algorithm finds near-optimal solutions