

MORA Routing and Capacity Building in Disruption-Tolerant Networks

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- mobile ad hoc wireless network
- end-to-end routing
- depends on the number of participants, their storage capacity, communication and movement patterns
- paper: focus is on performance

- improve performance by using MORA(=Multi-Objective Robotic Assistance)

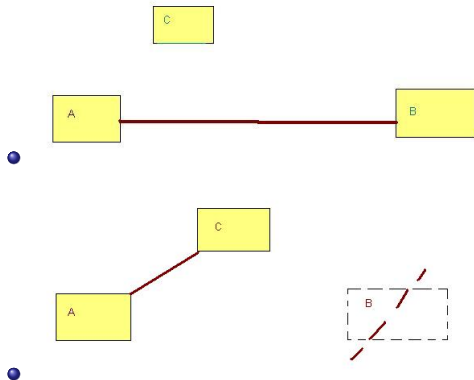
- autonomus agents to balance the network demand

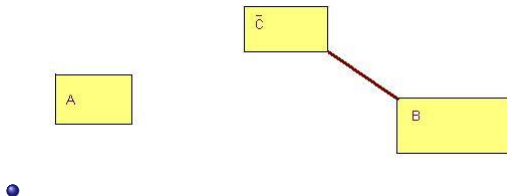


- adapt the motion of the agents to bandwidth and latency

- end to end communication
- mobile nodes
- nodes store and pass

Communication





routing algorithm "Drop-Least Encountered"

- estimate the distance by Likelihood
- peer = 0
- exchange by meeting
- peer=1 and send
- values sink by time and increase by frequent meetings

- NP hard
- reduce "dial-a-ride problem" a Traveling-Salesman-Problem
- a graph and the peers present participants
- message from A to B
- agent transfer if there is no communication

to get optimized motion in praxis

- multi-objective control methods
- generate complex behavior by put together simple
- controller present the behavoir with a control function

Subsumption:

- include one or more controllers per layer
- status information and can give commands
- Higher level controller can grasp status information of lower controller and overwrite their commands

Nullspace:

- concurrent execution of multiple controllers
- hierarchy
- lower-level controller have no effect of higher-level
- consists of all vectors x so that $Ax=0$
- collection of control commands

Nullspace example

- point $p=(x/y)$
- controller ϕ_1 changing state to achieve $x=0$
- controller ϕ_2 changing state to achieve $y=0$
- controller ϕ_2 orthogonal
- controller ϕ_2 operates in nullspace of controller ϕ_1

At each level of hierarchy the controller optimizes its action in nullspace of higher controllers

- Bandwidth
- Unique Bandwidth
- Message Latency
- Peer Latency

Bandwidth:

- total number of messages at a point in time
- schedule motion of agents to increase bandwidth
- space of transfer message should be used effectively
- agents have to consider their travel time for delivery

Unique Bandwidth:

- total number of unique messages which are currently active
- bandwidth should be used effectively
- agents should transmit messages not already in transmit

Message Latency:

- average amount for a message delivery
- for reducing agents should move towards peers which are sending or receiving many messages

Peer Latency:

- average time since a peer was last visited by agent
- participants should be visited intermittently

Distributed Network State Maintenance

- real network no global information available
- each peer has information about their own state
- all participants should have synchronized clocks
- each agent have information about participants
- two agents meet and exchange packages and routing protocol

Movement Controllers for Performance Metrics

- Total Bandwidth Controller ϕ T
- Unique Bandwidth Controller ϕ U
- Message Latency Controller ϕ D
- Peer Latency Controller ϕ p

Total Bandwidth Controller ϕ T:

- the chosen peer has the largest number of messages not seen by agent

Unique Bandwidth Controller ϕ U:

- choose peer which have the largest number of messages not present

Message Latency Controller ϕ D:

- choose peer whose average delivery time is the largest

Peer Latency Controller ϕ p:

- choose location n_i (last visited)
- Δt_{n_i} time since n_i has been visited
- $n = \operatorname{argmin}_{n_i} \Delta t_{n_i} + t_{n_i}$

Movement Controllers for Performance Metrics

- these controllers consider only a single performance metric
- generate conflicts for agents
- for example
 - two peers with a large number of messages
 - Total Bandwidth Controller ϕT only maximize these peers
 - other peers wouldn't be visited
 - not possible to combine bandwidth and latency controller
 - solution: coordinate controllers and select parameters

hierarchy

- $\phi_p < \phi_D < \phi_U < \phi_T$
- $\phi_i < \phi_j$ means j more important than i
- example Subsumption:

example Subsumption:

- Controller with highest priority decides
- lower priority get only a chance if all higher optimized their metric
- Controller ϕT optimize its bandwidth and give the next lower controller ϕU a chance
- if ϕU violated the total bandwidth ϕT get the control back

Comparison of Nullspace and Subsumption

- nullspace controller execute concurrently and subsumption at a time
- low nullspace controller can't influence higher controller

based on

- DTN
- 30 buses
- GPS devices on each bus
- location of each bus available

trace generator

- use GPS to reconstruct movement and timing of 9 buses routes in area of 8.24 km
- each route is presented by trace days
- advantage: buses can be varied

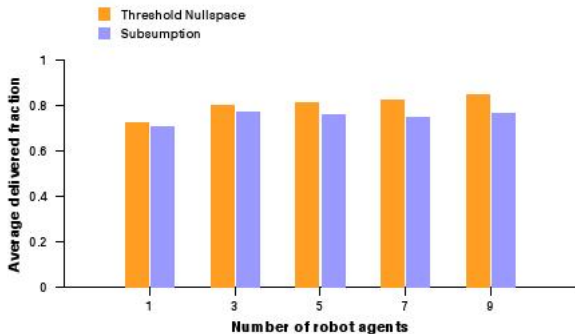
default values for each simulation:

- unlimited buffer
- load of 36 pkts/hour
- 10-Kbyte packets
- 9 buses, each on distinct routes
- agents increase the buffersource, so random choice 3 buses without free storage capacity

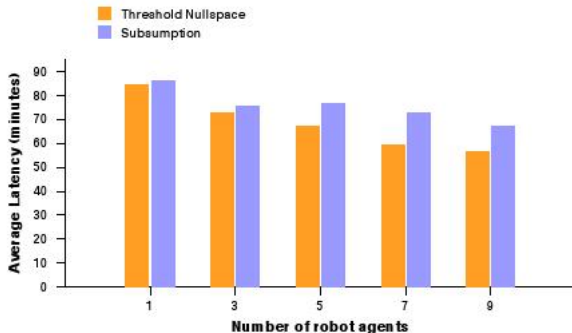
Protocols:

- Random: first packets to B and the rest by chance
- MaxProp: Each packet get a number, the packet with the lowest number is send first
- ME/DLE : packets get priority, lowest first
- FIFO : First in First out

Comparing Supsumption and Nullspace:

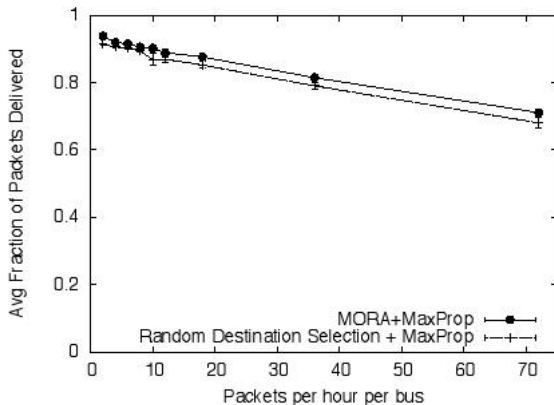


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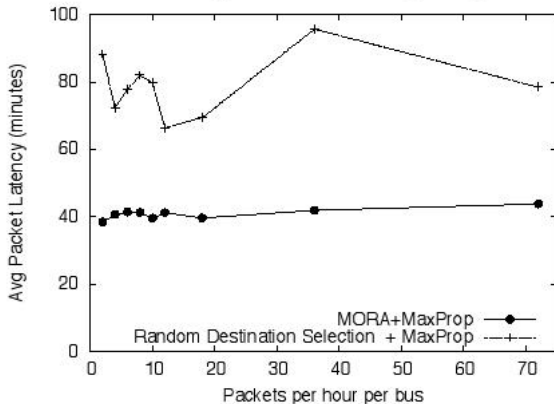


Experimental Evaluation

Comparison of MORA to movement using random selection

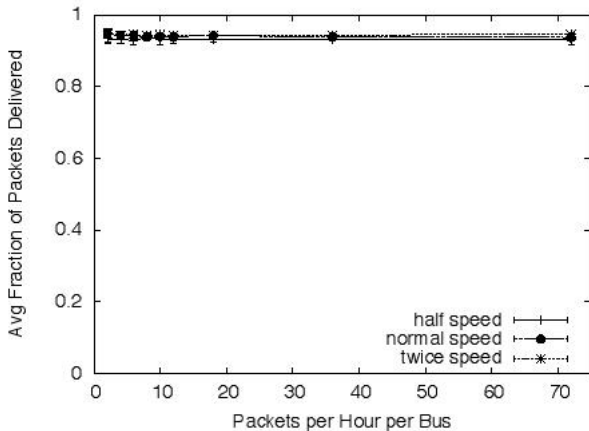


Comparison of MORA to movement using random selection



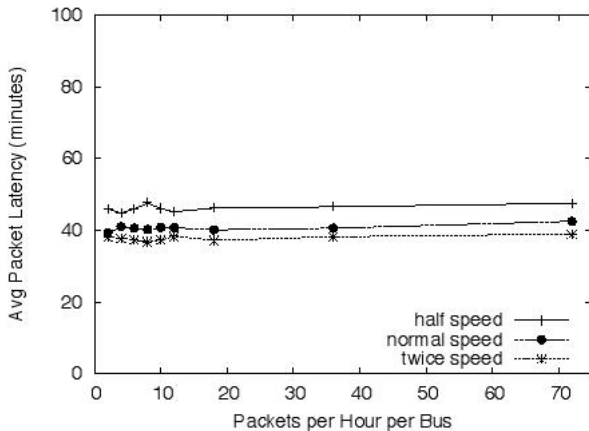
Experimental Evaluation

Comparison of MORA+MaxProp by varied speeds of agents



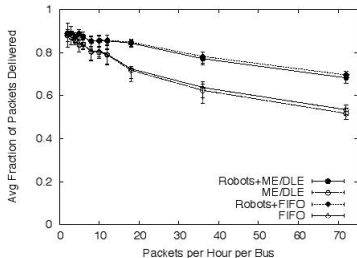
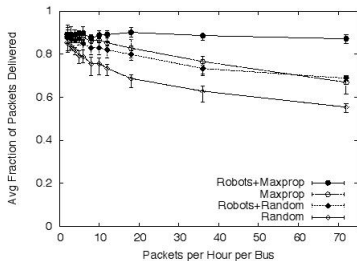
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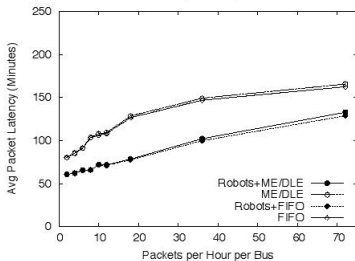
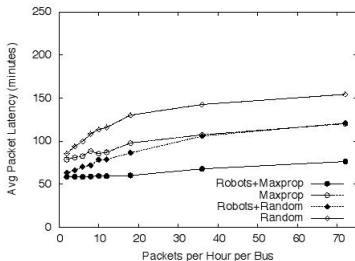
Experimental Evaluation

packet delivery rate as network load increases



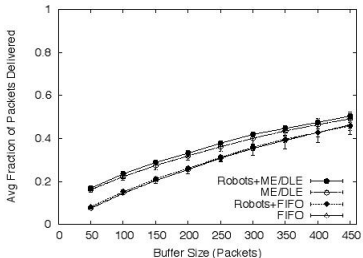
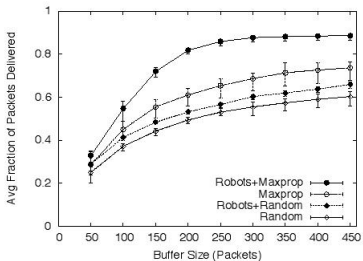
Experimental Evaluation

packet latency as network load increases



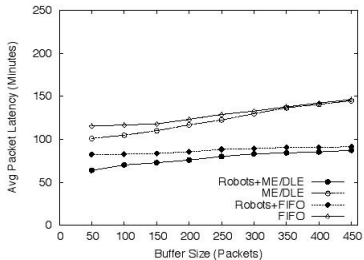
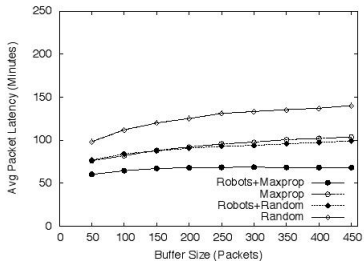
Experimental Evaluation

packet delivery rate as local buffer size increase



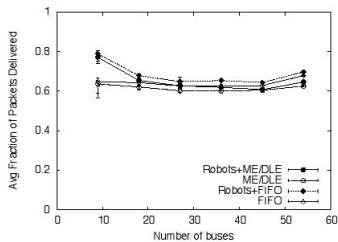
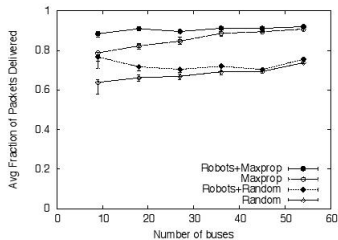
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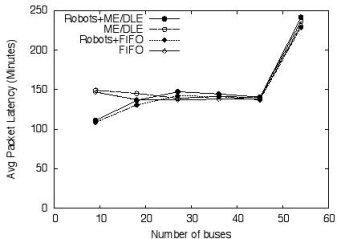
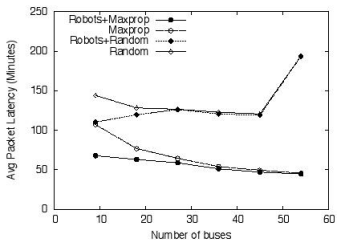
Experimental Evaluation

packet delivery rate with an increasing number of buses



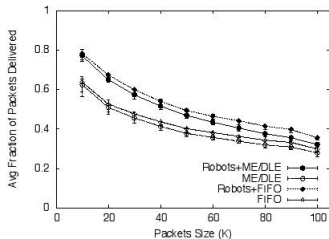
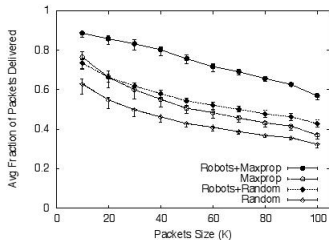
Experimental Evaluation

packet latency with an increasing number of buses



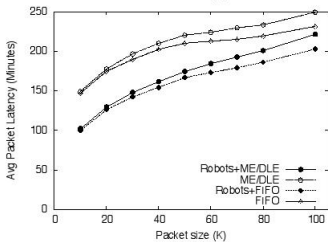
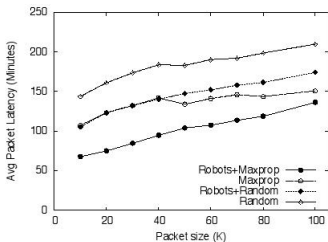
Experimental Evaluation

packet delivery rate as packet size increases



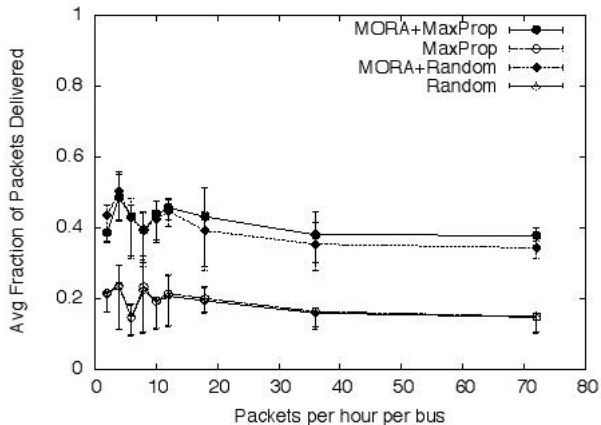
Experimental Evaluation

packet latency as packet size increases



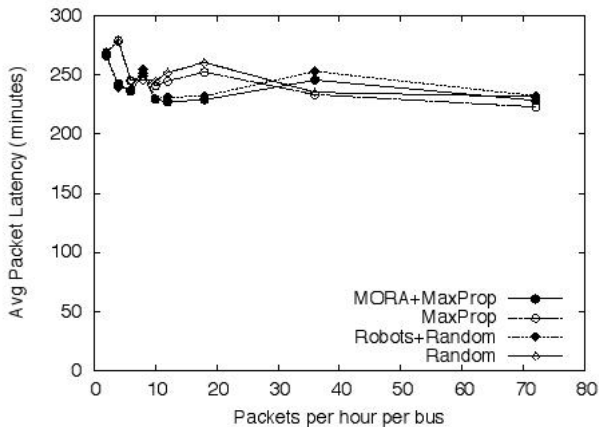
Experimental Evaluation

Synthetic mobility model: delivery rate as packet load increases



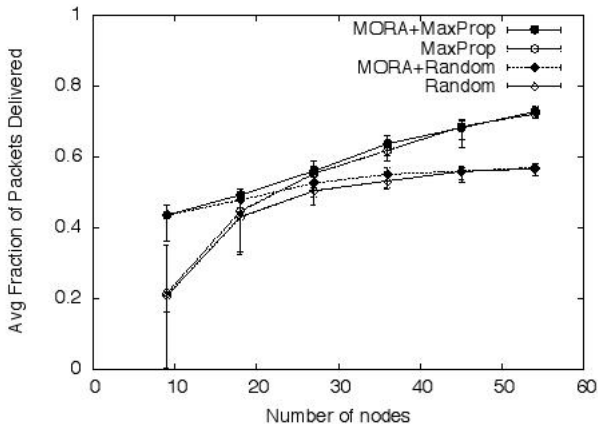
Experimental Evaluation

Synthetic mobility model: latency as packet load increases



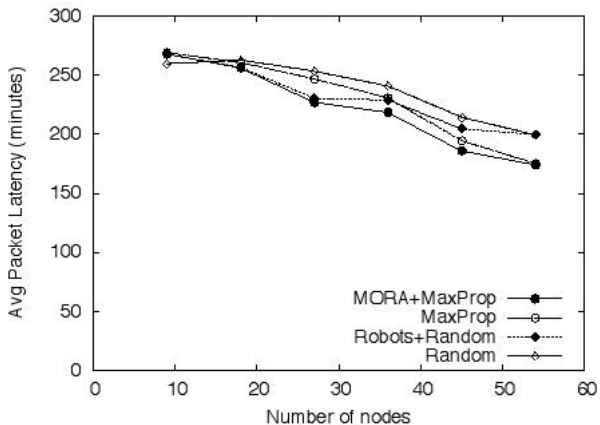
Experimental Evaluation

Synthetic mobility model: delivery rate as number of nodes increases



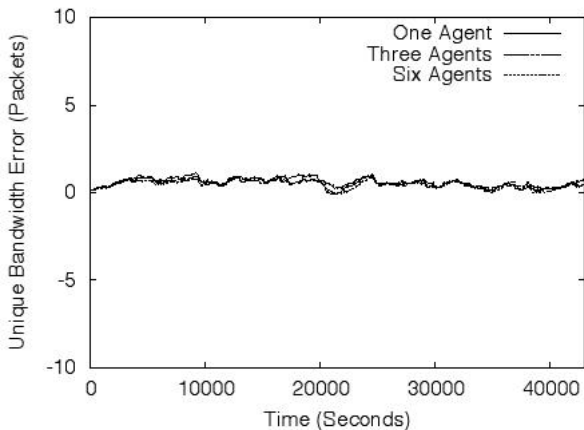
Experimental Evaluation

Synthetic mobility model: latency as number of nodes inceases



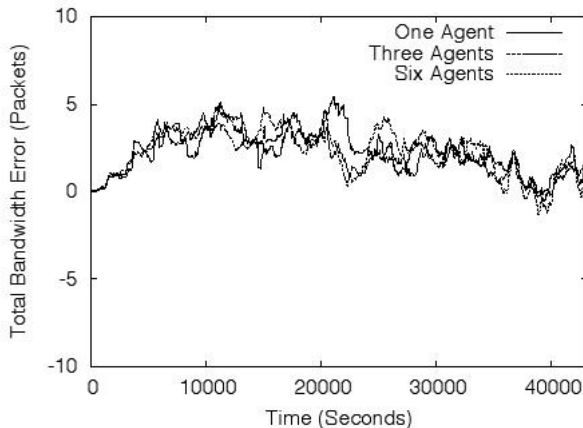
Experimental Evaluation

Bandwidth error



(a) Bandwidth Error

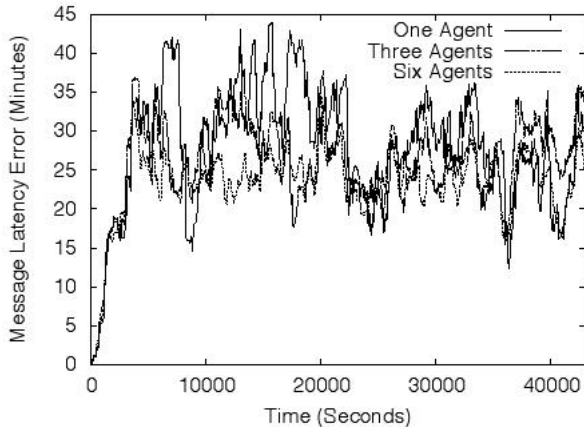
Bandwidth error



(b) Bandwidth Error

Experimental Evaluation

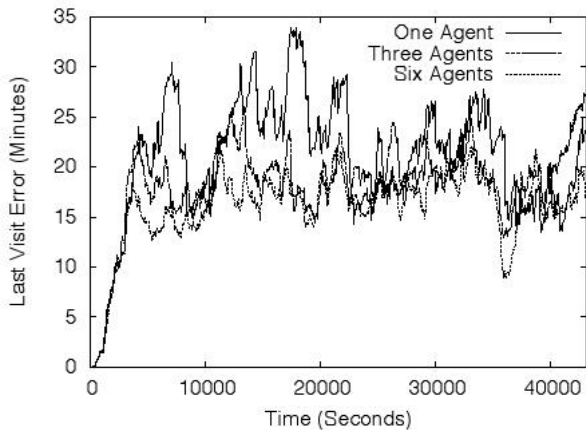
latency error



(c) Latency Error

Experimental Evaluation

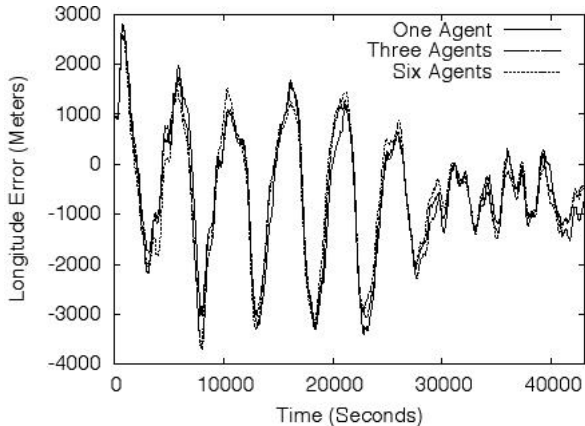
latency error



(d) Last Visited Error

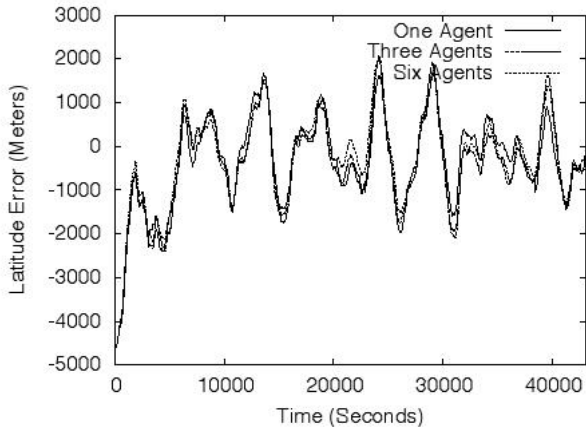
Experimental Evaluation

Agent longitude error



(e) Agent Longitude Error

Agent longitude error



(f) Agent Latitude Error

- To improve the performance of DTN use MORA
- MORA sets autonomous agents in the DTN
- To get the optimize movement use control methods

Thank you for your attention