

7 R, ZPP, BPP, Savitch and ATMs

R TM - accepts if
 50% of comp. paths
 are accepting
 - rejects if all leaves
 are rejecting.



$$R = RP = RTIME(n^{O(n)})$$

$$ZPP := R \cap co-R$$

$$PRIME := \{bin(m) \mid m \text{ is a prime number}\}$$

$$PRIME \in ZPP$$

$$P \subseteq ZPP, PRIME \in P!$$

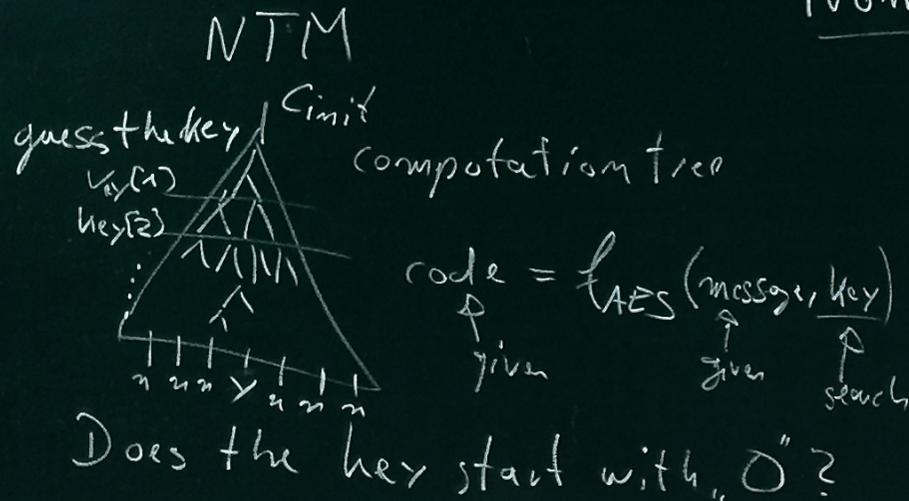
BPP: bounded pol. time
 accepts if more than $\frac{2}{3}$ accept
 rejects if more than $\frac{2}{3}$ reject

$$R \cup co-R \in BPP \subseteq PP$$

if $NP \subseteq BPP \Rightarrow R = NP!$

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Non-Determinism



- Unambiguous TM (UTM)
there is at most one accepting leaf in the computation tree

UTM accepts, if it exists

$$UP := UTIME(n^{O(n)})$$

$$P \subseteq UP \subseteq NP$$

$$\begin{matrix} \text{in} \\ \oplus P \end{matrix}$$

- Few-TM
 promise: the number of
 accepting leafs is bounded
 by a polynomial.

$$\text{FewTIME}(n^{O(n)}) =: \overline{\text{FewP}}$$

$$\text{UP} \subseteq \overline{\text{FewP}} \subseteq \text{NP}$$

$$\overline{\text{FewP}} \subseteq \oplus \text{P} \quad !$$

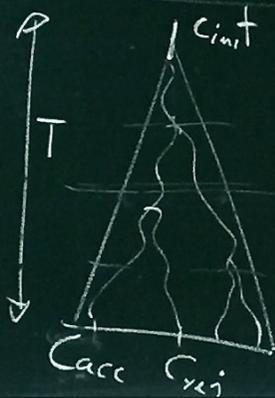
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Theorem of Savitch

Theorem For any space-constructible function $s: \mathbb{N} \rightarrow \mathbb{N}$, $s(n) \geq n$

$$NSPACE(s(n)) \subseteq SPACE(s(n)^2)$$

Proof:



computation tree
 unique accepting
 configuration c_{acc}
 time $T \in 2^{c \cdot s(n)}$
 $c > 0$ constant

Configuration

$$Q \times |\Sigma|^{k \cdot s(n)} \times \{\perp, \triangleright\}^k$$

$c \cdot s(n)$

$$\# \text{config} \leq 2$$

1. BAD idea: Backtrack in
in the comp. tree

Size of the stack:

- one entry config Space: $s(n)$

- T entries

$$\Rightarrow \text{Size } s(n) \cdot 2^{s(n)}$$

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$$\text{CanYield}(C_1, C_2, 0) = (C_1 = C_2) \quad \checkmark$$

$$\text{CanYield}(C_1, C_2, 1) = \begin{cases} 1, & C_1 \neq C_2 \\ 0, & \text{else} \end{cases} \quad \checkmark$$

$$\text{CanYield}(C_1, C_2, t) =$$

\exists conf. C s.t.

$$\text{CanYield}(C_1, C, \lfloor \frac{t}{2} \rfloor) \wedge \text{CanYield}(C, C_2, \lceil \frac{t}{2} \rceil)$$

$$t = \lfloor \frac{t}{2} \rfloor + \lceil \frac{t}{2} \rceil$$

$$5 = \lfloor 2.5 \rfloor + \lceil 2.5 \rceil = 2 + 3 = 5$$

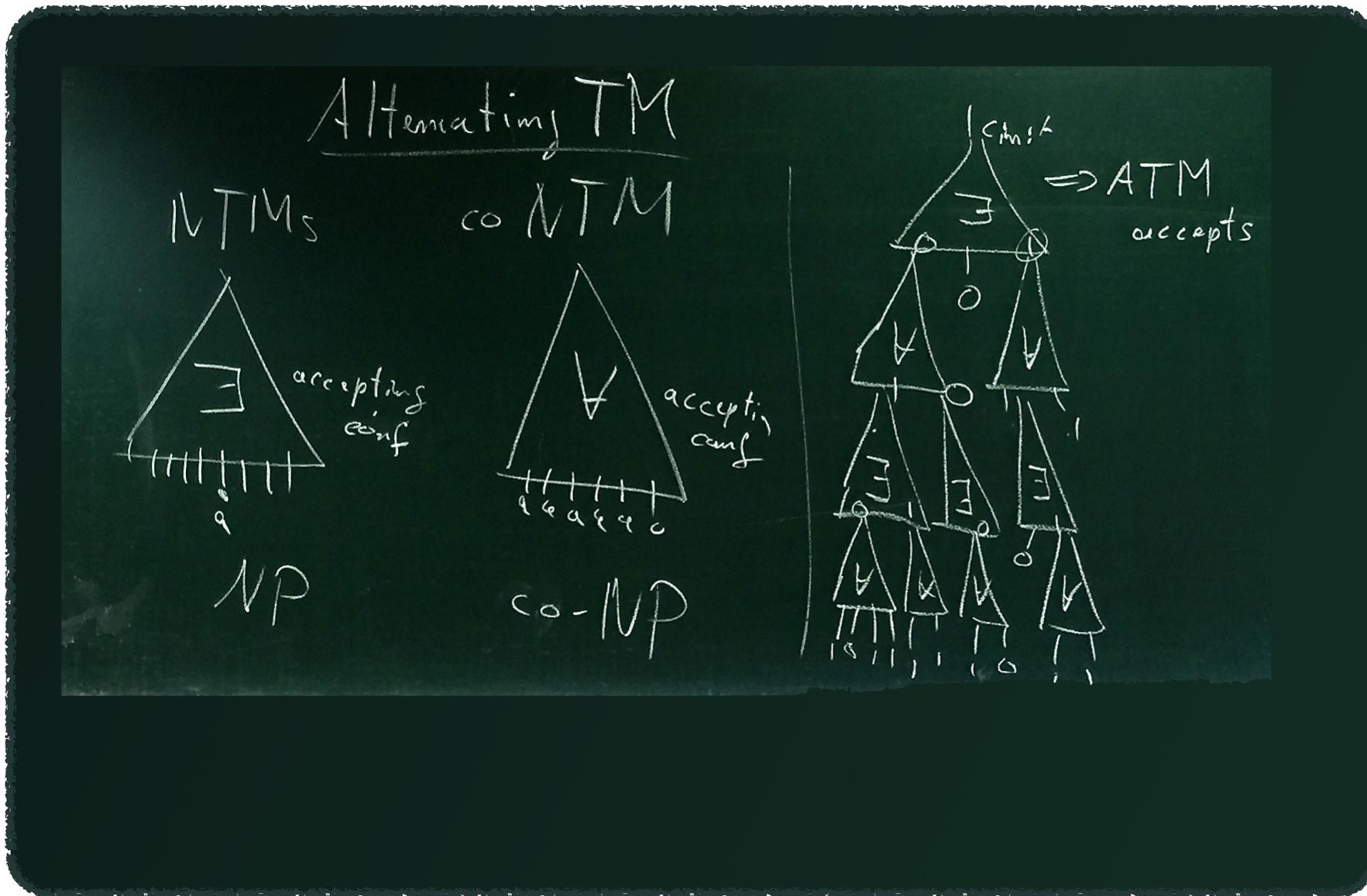
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Alg CanYield(C_1, C_2, ϵ)
 if $t=0$ then if $C_1 = C_2$ then
 RETURN true
 else RETURN false
 else if $t=1$ then if $C_1 \neq C_2$ then
 else true
 else false
 else forall c_1, c_2 of size $s(n)$ do
 if CanYield($C_1, C_2, \frac{\epsilon}{2}$)
 and CanYield($C_1, C_2, \frac{\epsilon}{2}$)
 then RETURN TRUE
 od
 RETURN FALSE

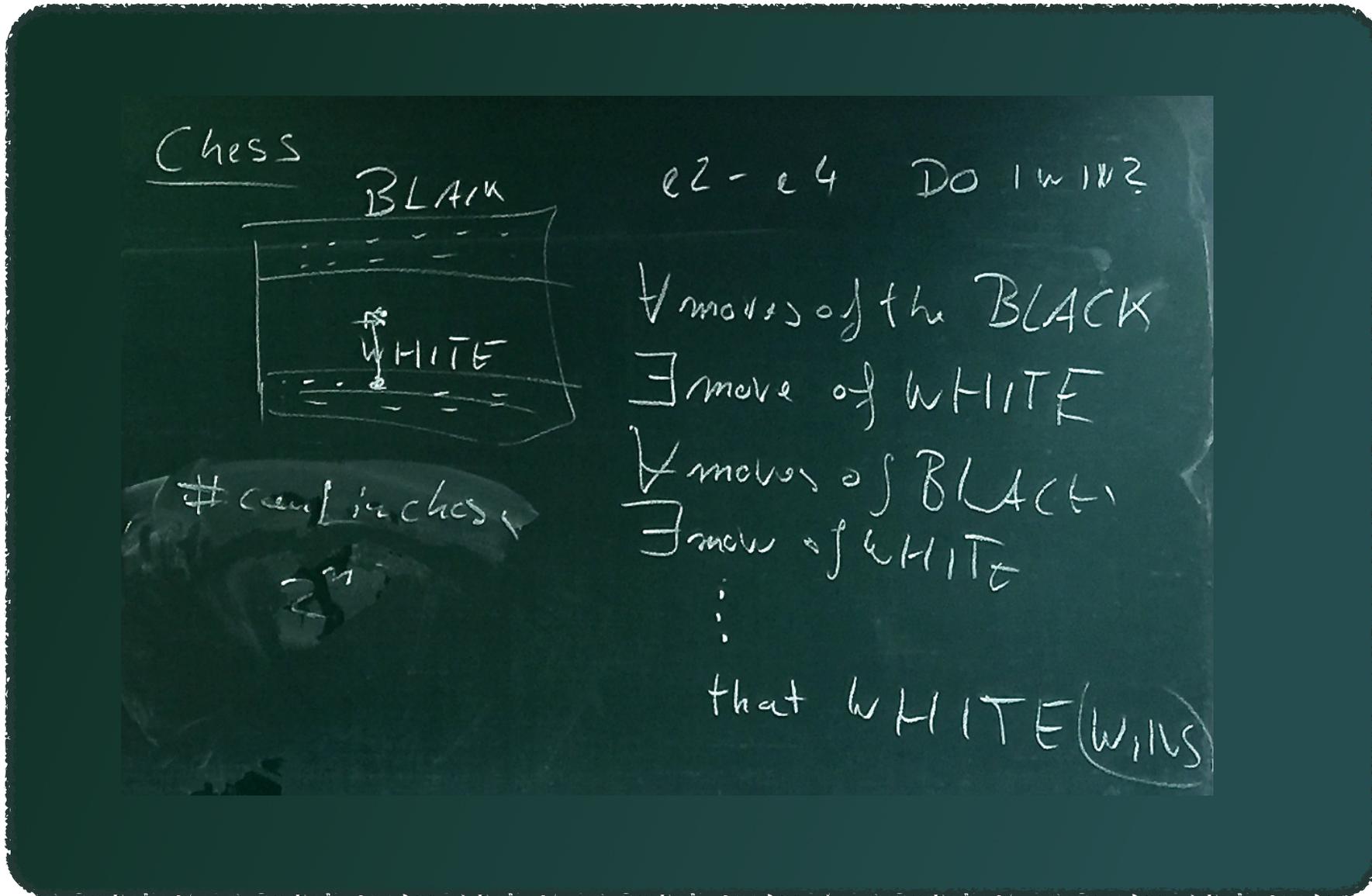
space needed per
 recursion : $O(s(n))$
 stack size $\log T + 1$
 $T \rightarrow T/2 \rightarrow T/4 \rightarrow \dots 1$
 $\log T = O(s(n))$
 overall space needed
 is $O(s(n)^2)$
 Now compute
 CanYield(C_1, C_2, ϵ)

\Rightarrow
 PSPACE = NPSPACE

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ATM has states Q

$$Q = \{q_{acc}\} \cup \{q_{rej}\} \cup Q_{\exists} \cup Q_{\forall}$$

ATM M accepts in conf C .

1. C is in q_{acc}
2. C is in Q_{\exists} and $\exists C' : C \vdash C'$ and C' accepts
3. C is in Q_{\forall} and $\forall C' : C \vdash C' \Rightarrow C'$ accepts

$$AP = ATIME(n^{O(n)})$$

$$AP = PSPACE$$

(next time...)