Freiburg, the 15th November, 2011 Discussion 21th November, 2011

## Exercise for the lecture

## **Distributed Storage and Computer Forensic**

## Winter 2011/12

Sheet 4

## **EXERCISE 4:**

Consider the following (2, 4)-B-Tree.



- 1. Insert the elements 80, 75, 65 and 70.
- 2. Delete 10, 15 and 30.

**Trickquestion** Consider a B-tree consisting of a root and  $\sqrt{n}$  children each containing  $\sqrt{n}$  elements. A binary search for an element would then cost  $O(\log(\sqrt{n}))$  (2 binary searches in  $\sqrt{n}$  elements). From other classes you know that binary search is bounded by  $\Omega(\log n)$ ! How is that possible?