



**Exercises**  
**Distributed Systemes: Part 2**  
**Summer Term 2014**  
3.7.2014

## 4. Exercise sheet: Refresh Concurrency Control and Recovery

### Exercise 1

Consider the following schedules.

$S_1$ :  $R_3X R_2Y W_2Y R_1Y W_1Y R_2X W_2X R_1X W_1X W_3Z$ .

$S_2$ :  $R_3X R_2Y W_2Y R_1Y W_1Y R_2X W_2X R_1X W_1X W_3Y$ .

$S_3$ :  $R_1Y W_1Y R_2Y W_2Y R_2X W_2X R_3Z W_3X R_1X W_1X$ .

For each schedule give its conflict graph. Which schedules are serializable, which are not?

### Exercise 2

Assume on a database three transactions are being executed.

- a) The transactions are of the form:
- $$\begin{array}{l} T_1 : RA \quad WA \\ T_2 : RA \quad WA \\ T_3 : RA \quad WA \end{array}$$

(i) How many serial schedules do exist for  $T_1, T_2, T_3$ ? Give the reasons!

(ii) How many serializable schedules do exist for  $T_1, T_2, T_3$ , which are not serial ones? Give the reasons!

- b) The transactions are of the form:
- $$\begin{array}{l} T_1 : RA \quad WC \\ T_2 : RB \quad WA \\ T_3 : RC \quad WD \end{array}$$

(i) How many schedules do exist for  $T_1, T_2, T_3$ , which are not serializable? Give the reasons!

(ii) Applying 2-phase-locking, is it possible that all serializable schedules of  $T_1, T_2, T_3$  may occur? Give the reasons!

### Exercise 3

Consider the schedule  $S$ :

$$\begin{array}{l} T_1 : \quad \quad \quad R(X) \quad \quad \quad W(Y) \\ T_2 : \quad \quad \quad \quad \quad R(Y) \quad \quad \quad W(Y) \\ T_3 : R(Z) \quad W(Y) \end{array}$$

(a) Demonstrate that  $S$  is not (conflict-) serializable.

(b) We call two schedules equivalent, whenever (i) they are built out of the same transactions, (ii) in both schedules the transactions read the same values, and (iii) both schedules produce the same final state of the database. Demonstrate that the serial schedule  $T_1 T_3 T_2$  and schedule  $S$  are equivalent.

