



ALBERT-LUDWIGS-  
UNIVERSITÄT FREIBURG

# Algorithms for Radio Networks

**Routing in MANET: Link Reversal, OLSR, ZRP**

University of Freiburg  
Technical Faculty  
Computer Networks and Telematics  
Christian Schindelhauer



# Link Reversal

- ▶ **Gafni, Bertsekas,**
  - *Distributed Algorithms for Generating Loop-Free Routes in Networks with Frequently Changing Topology,*  
IEEE Transactions on Communications, Vol. 29, No. 1  
pp. 11-18, IEEE, January 1981
- ▶ **Routing protocol**
  - with special repair mechanism

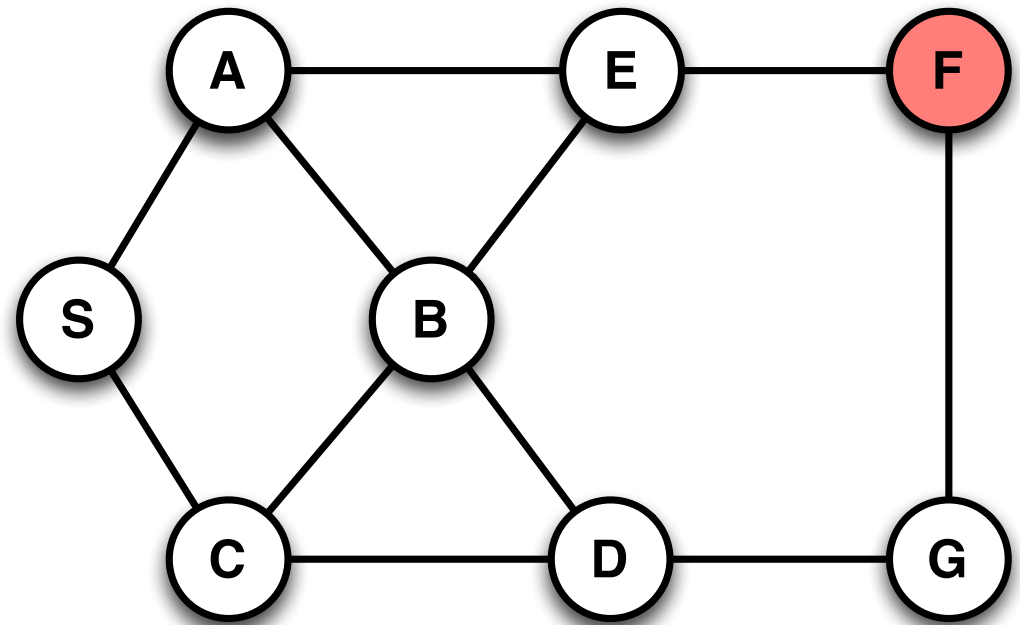
# Link Reversal

► **For each target node a direction for each edge is defined**

- all edges point towards the target
- e.g. by flooding and topological sorting

► **Routing**

- Pick any outgoing edge and send packet



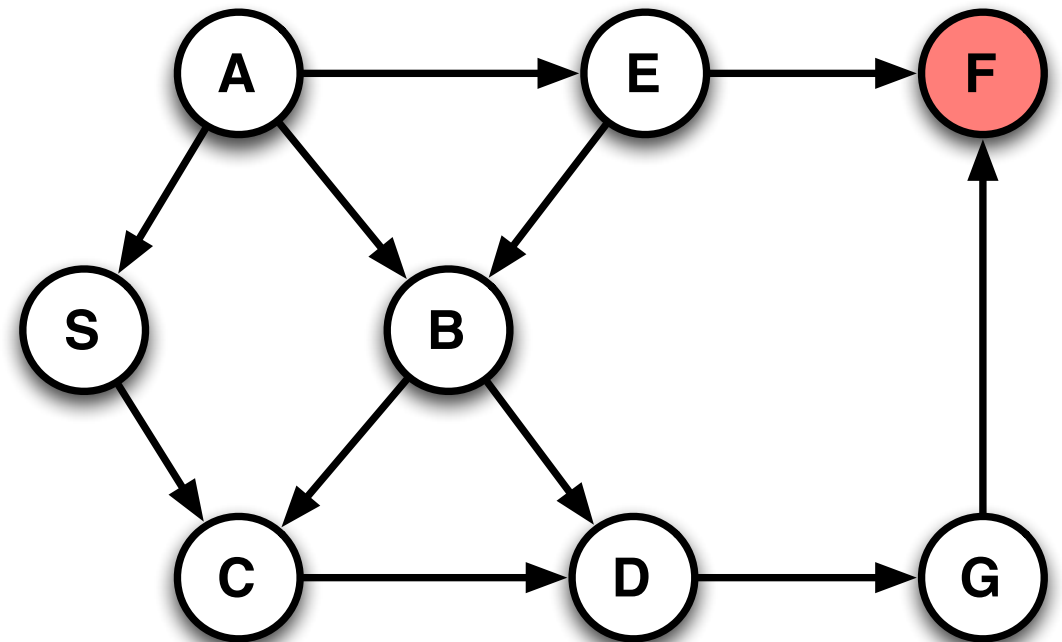
# Link Reversal

- ▶ **Directed acyclic graph (DAG) for each target**

- is preserved also in the case of failing edges

- ▶ **Connections are symmetrical**

- direction are only virtual

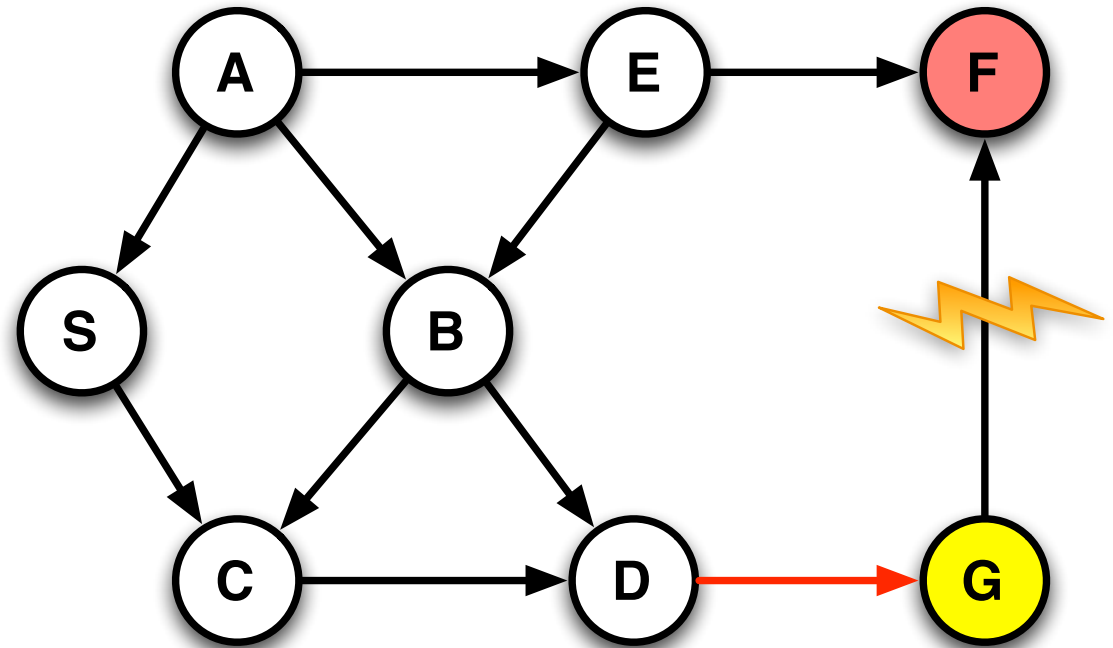


# Link Reversal

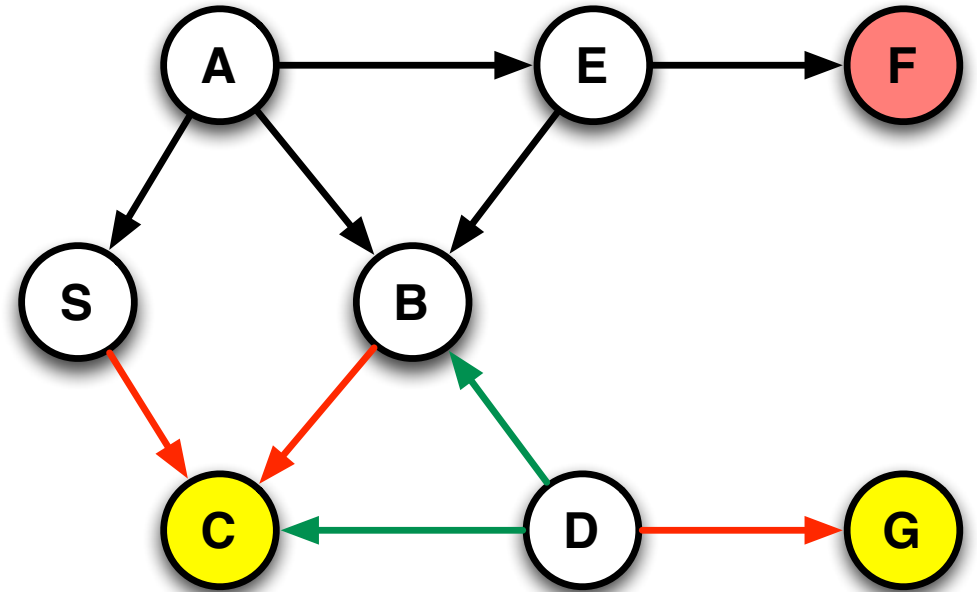
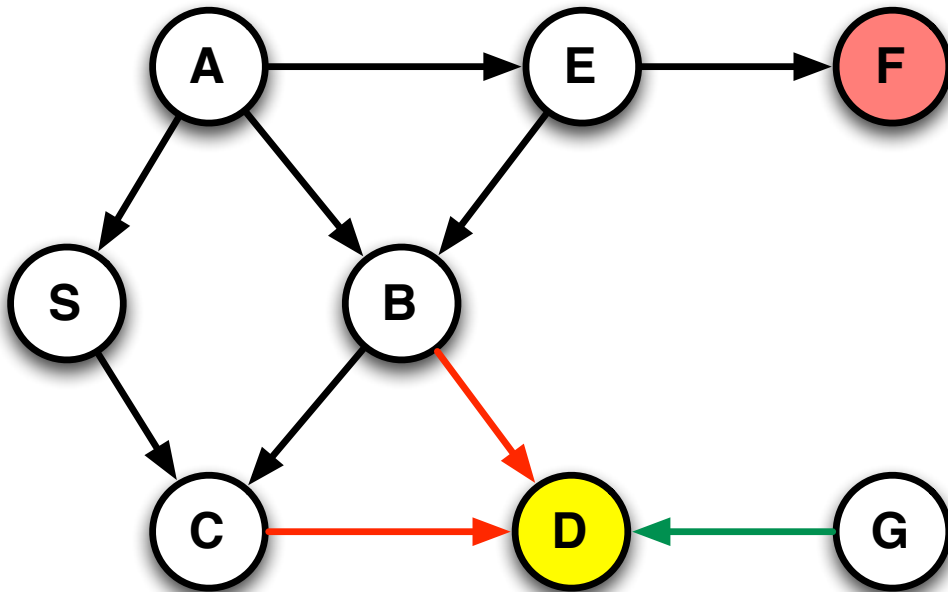
## ► Link F-G is lost

## ► Repair

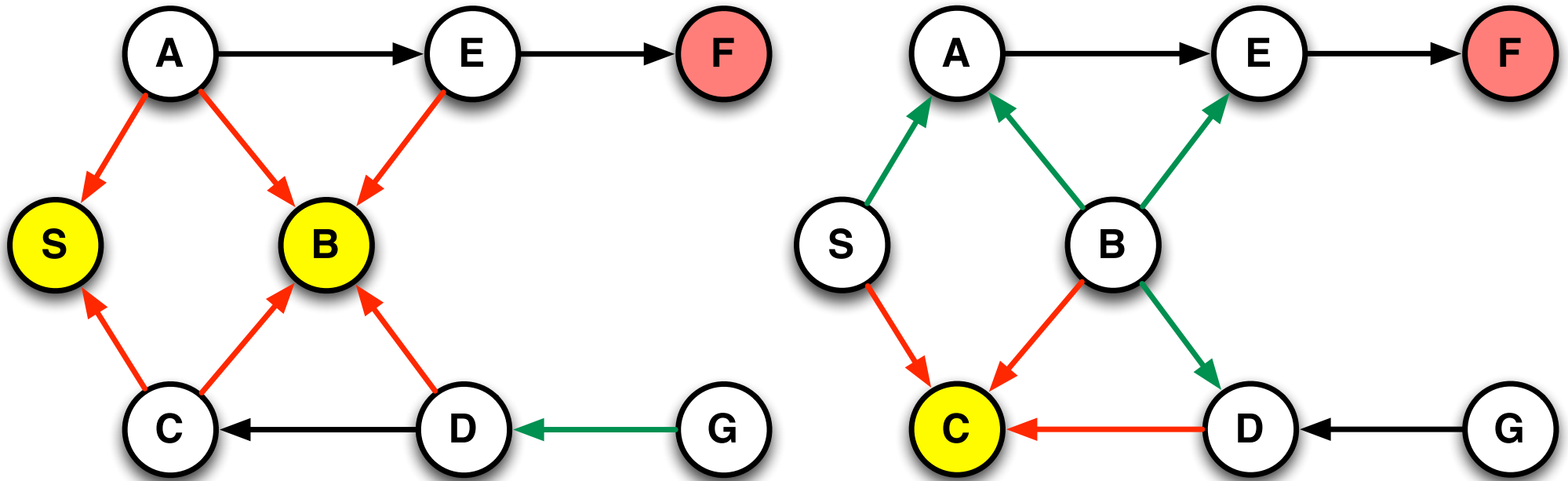
- All nodes without outgoing edges change the orientation of all incoming edges



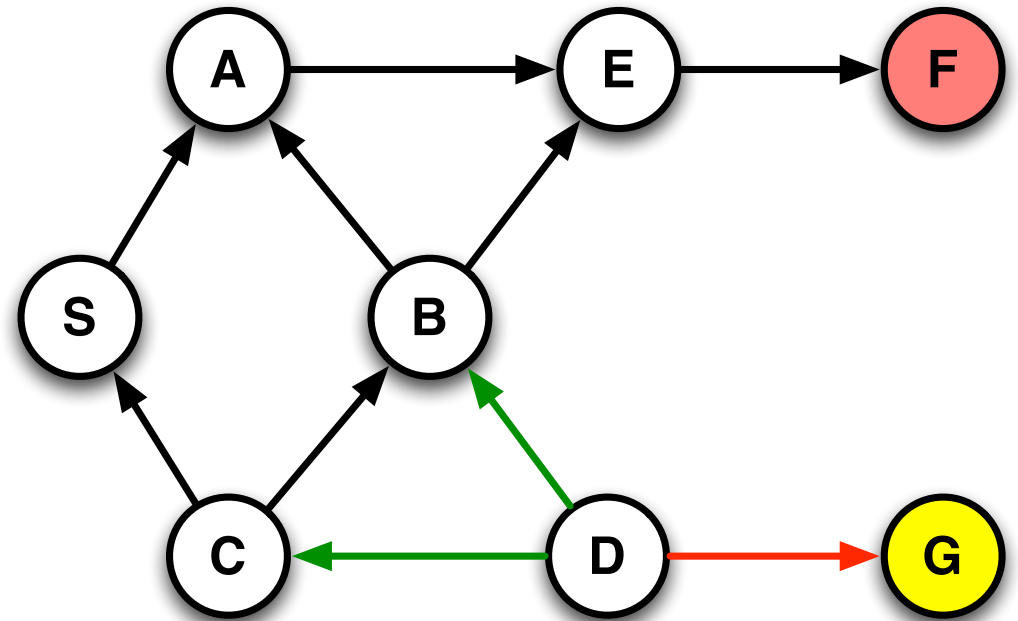
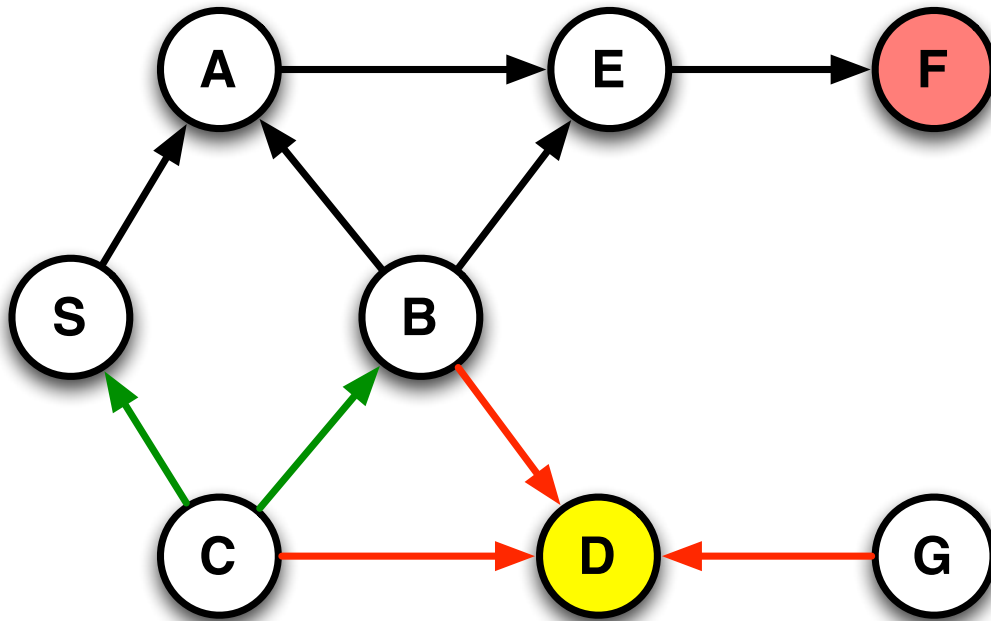
# Link Reversal



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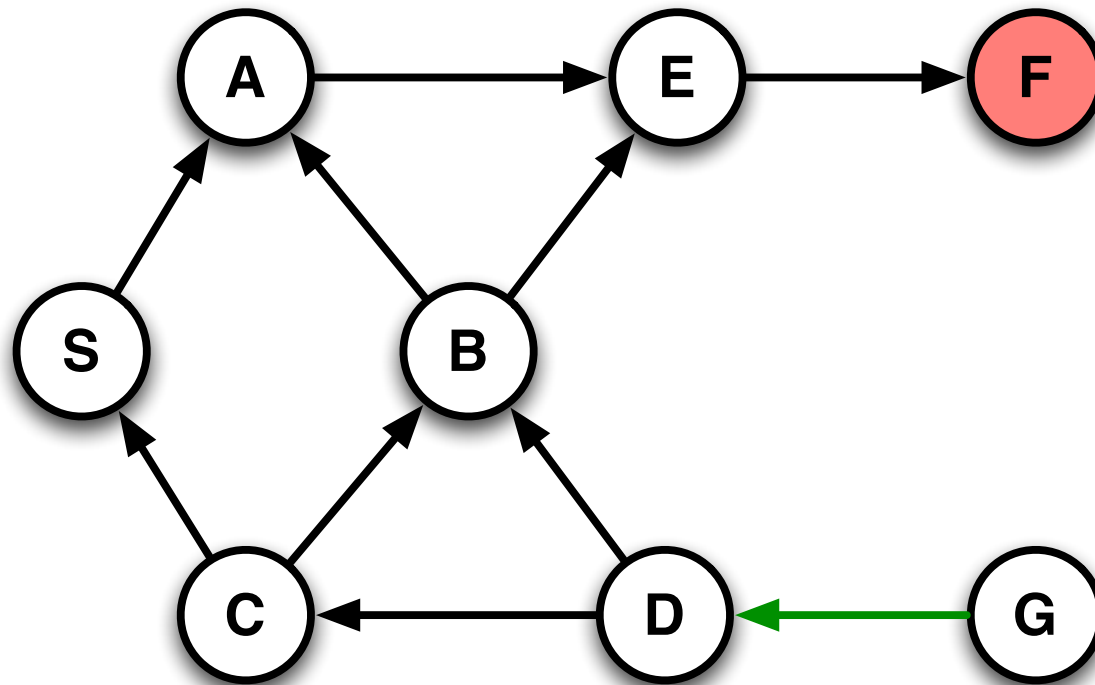


# Link Reversal





# Link Reversal



# Link Reversal

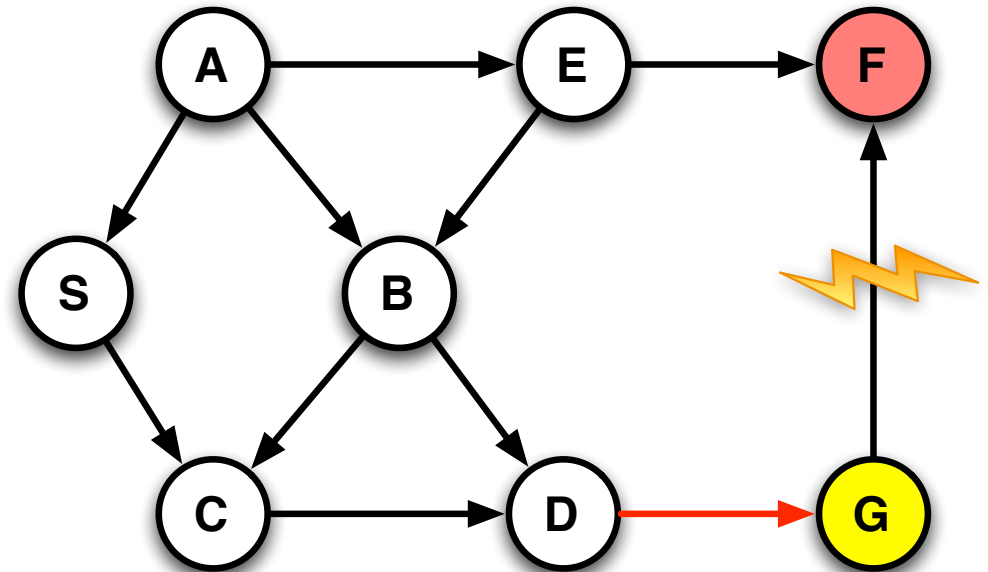
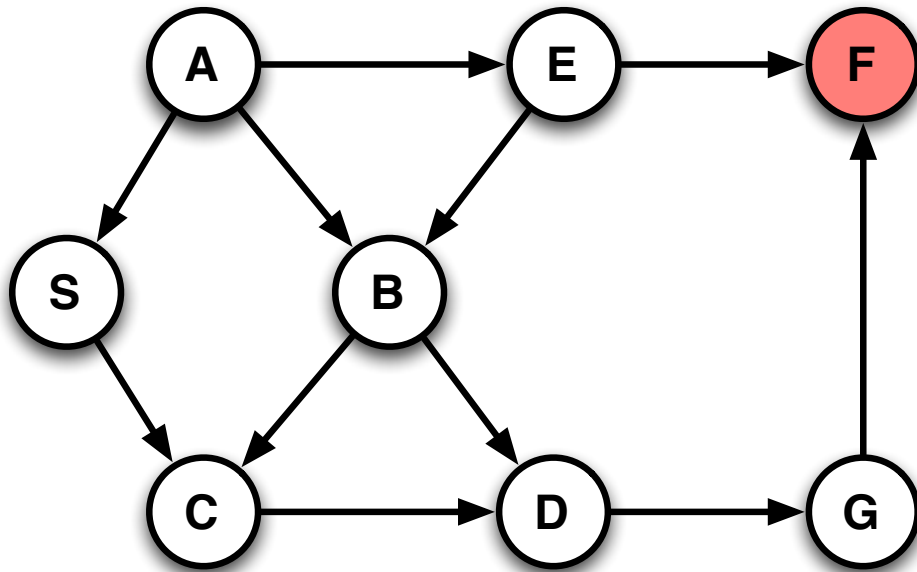
- **Motivation**
  - Link reversal should cause only local changes
  - Not necessarily the case
- **Repair is initiated,**
  - when the first packet is sent
- **Method known as a full reversal**

# Partial Reversal

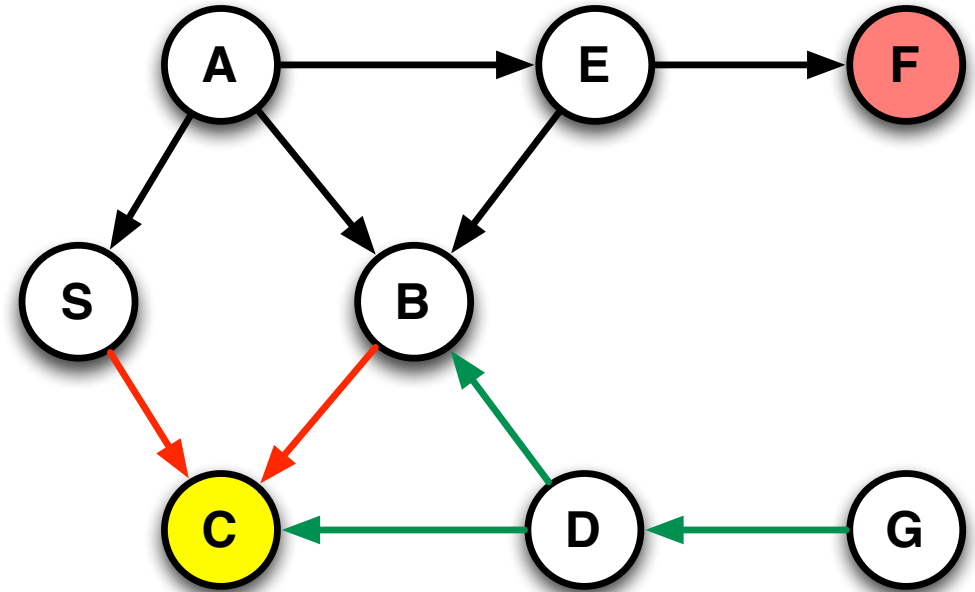
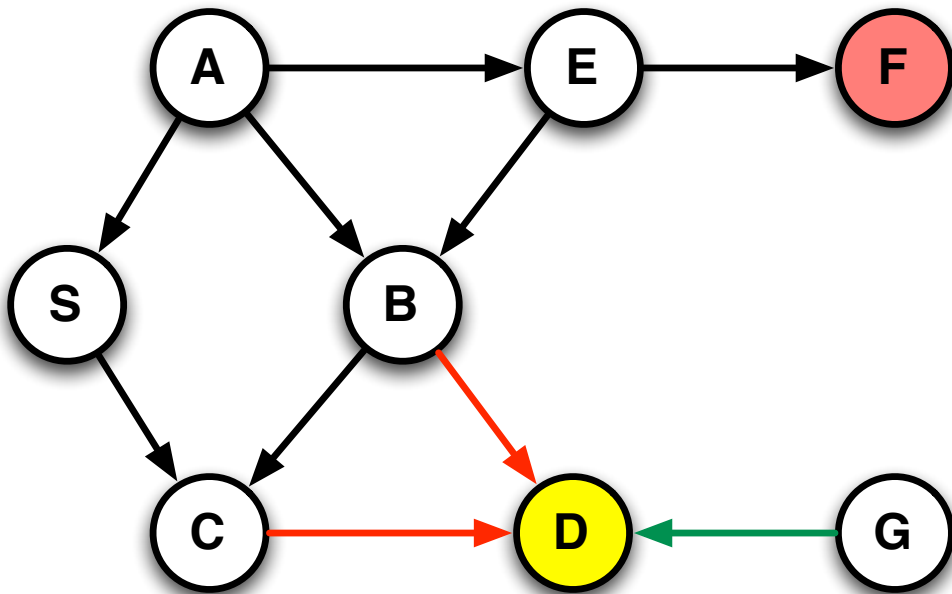
## ► Partial reversal

- Only the edges are reversed which have not been reversed before
- If all edges have been already reversed, then reverse all edges again

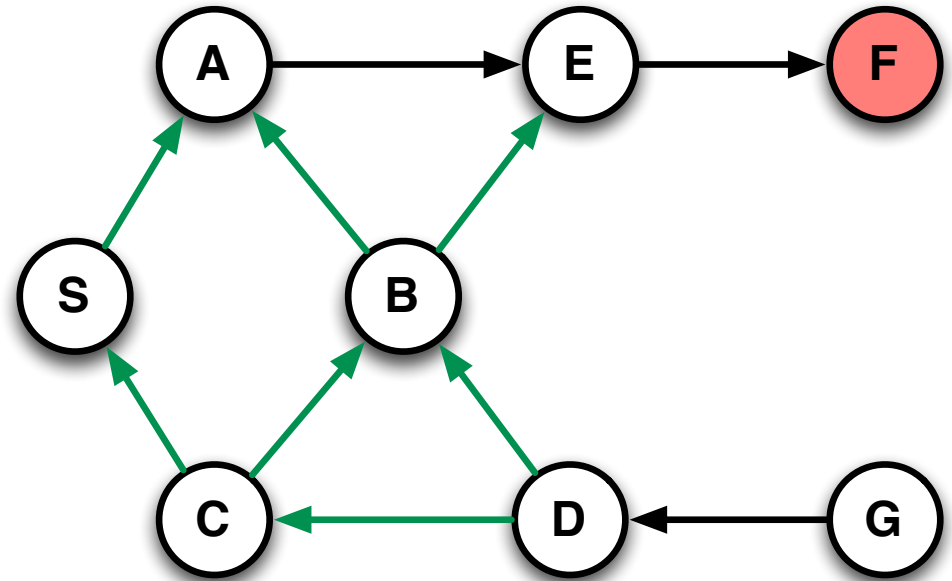
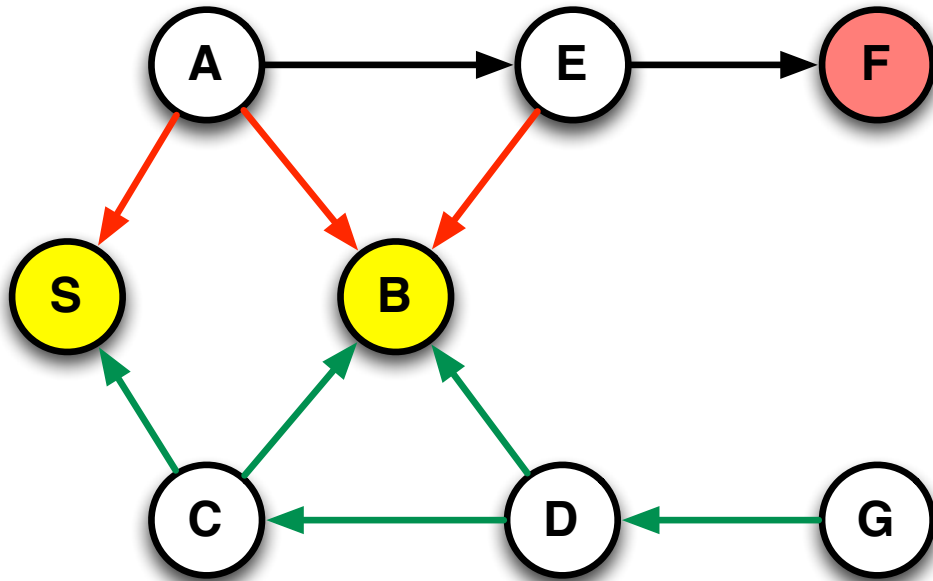
# Partial Reversal



# Partial Reversal



# Partial Reversal



# Link Reversal

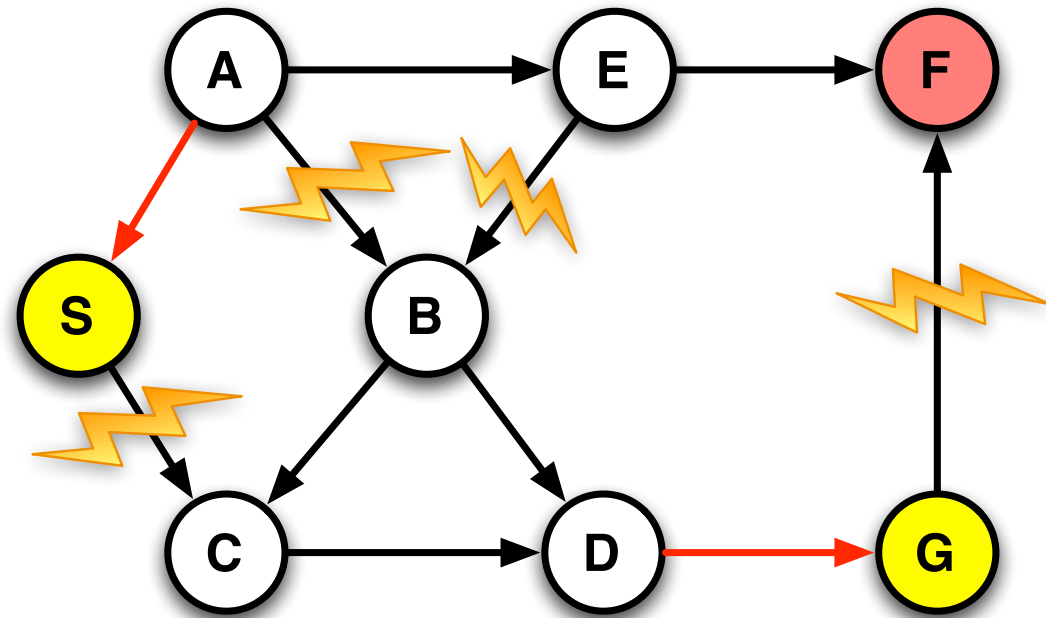
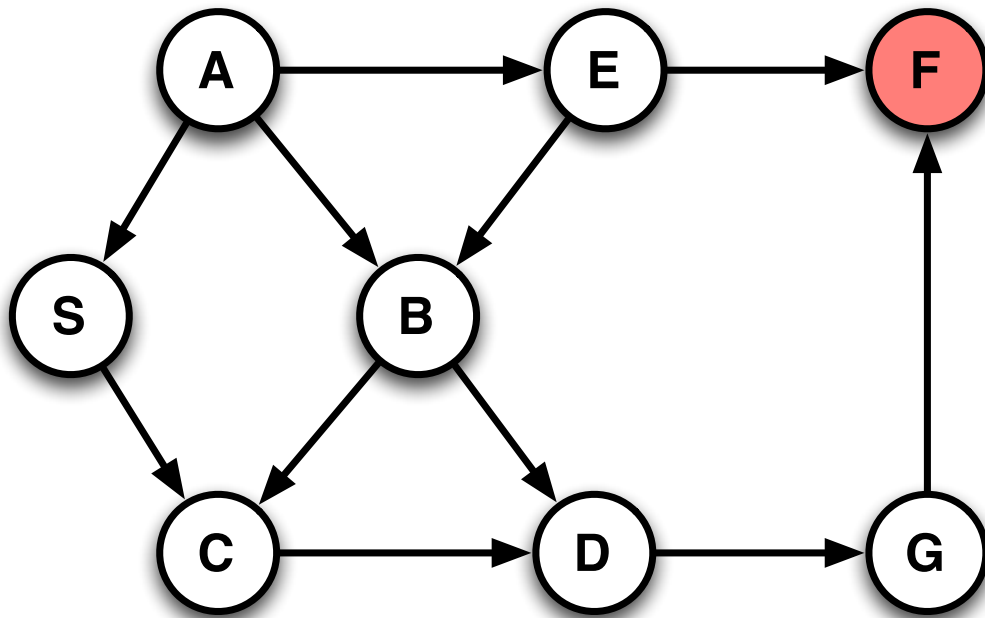
## ► Advantages

- Link reversal intends local repair
- Several substitute routes (potentially) available

## ► Disadvantages

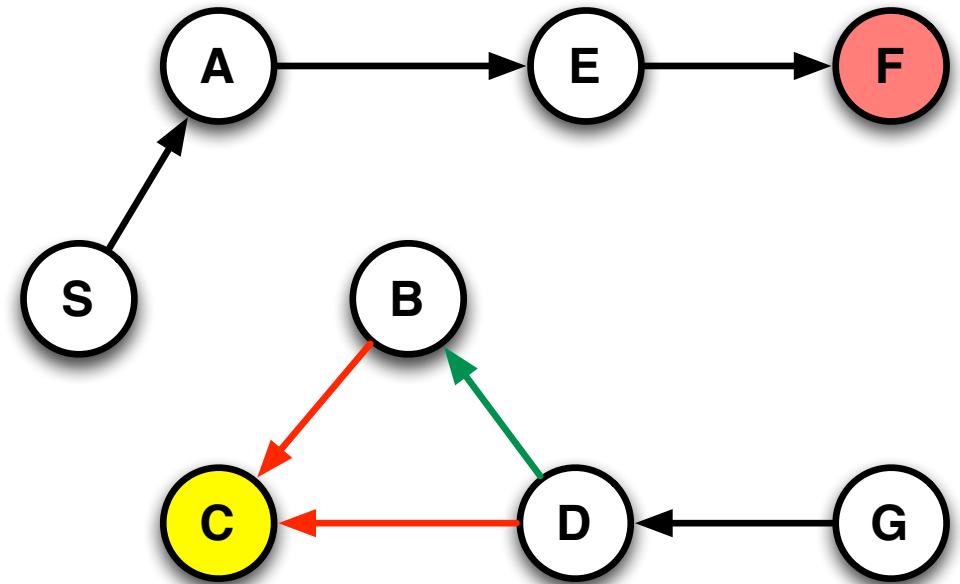
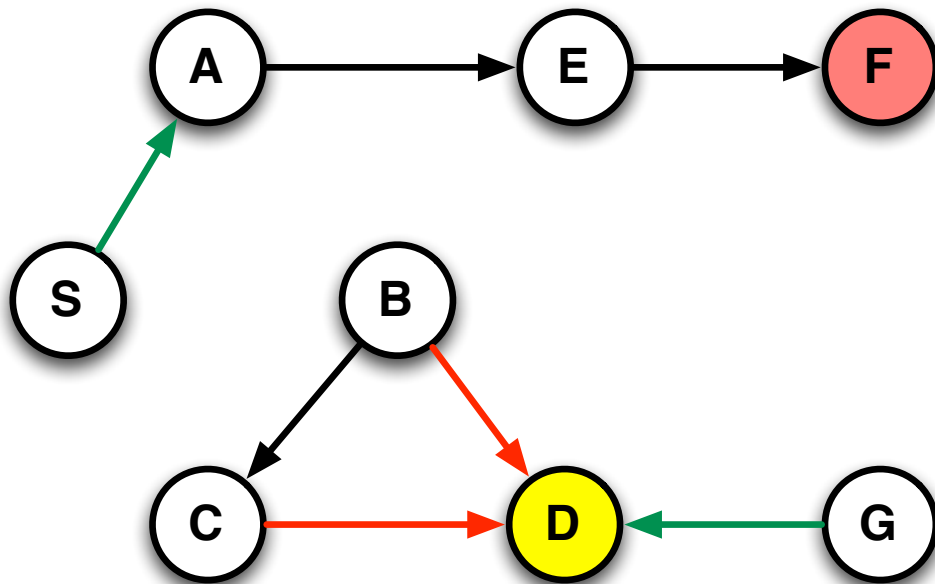
- Connection errors must be detected
  - Hello messages cause additional traffic
- If network is partitioned, the repair mechanism does not terminate

# Link Reversal if Network is Partitioned

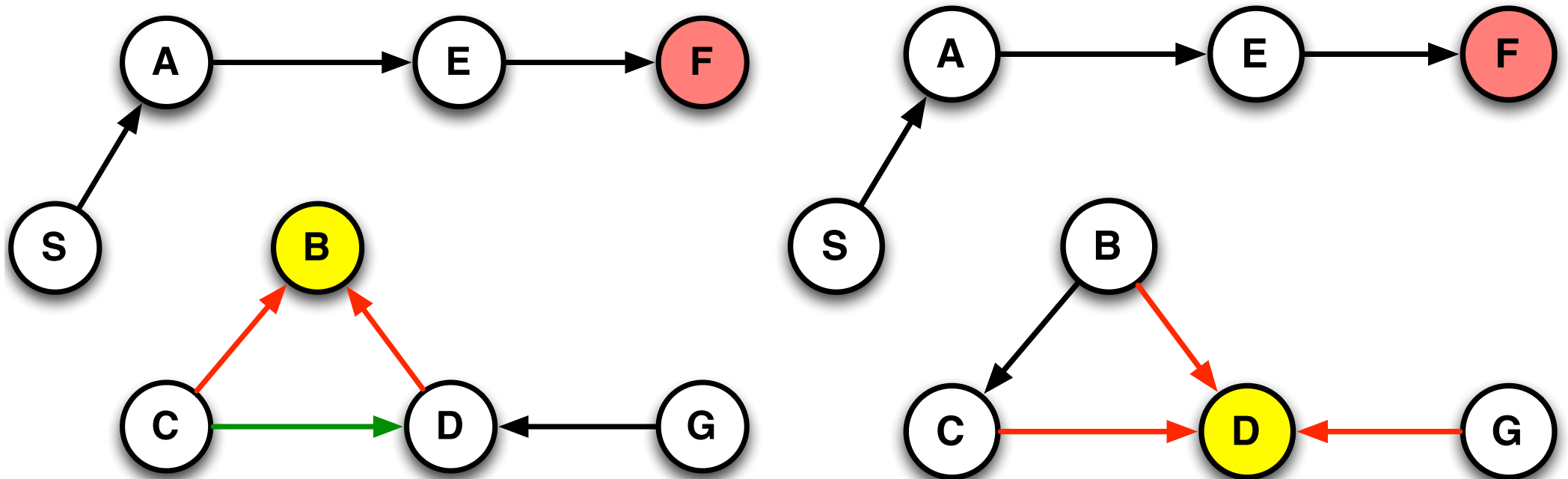




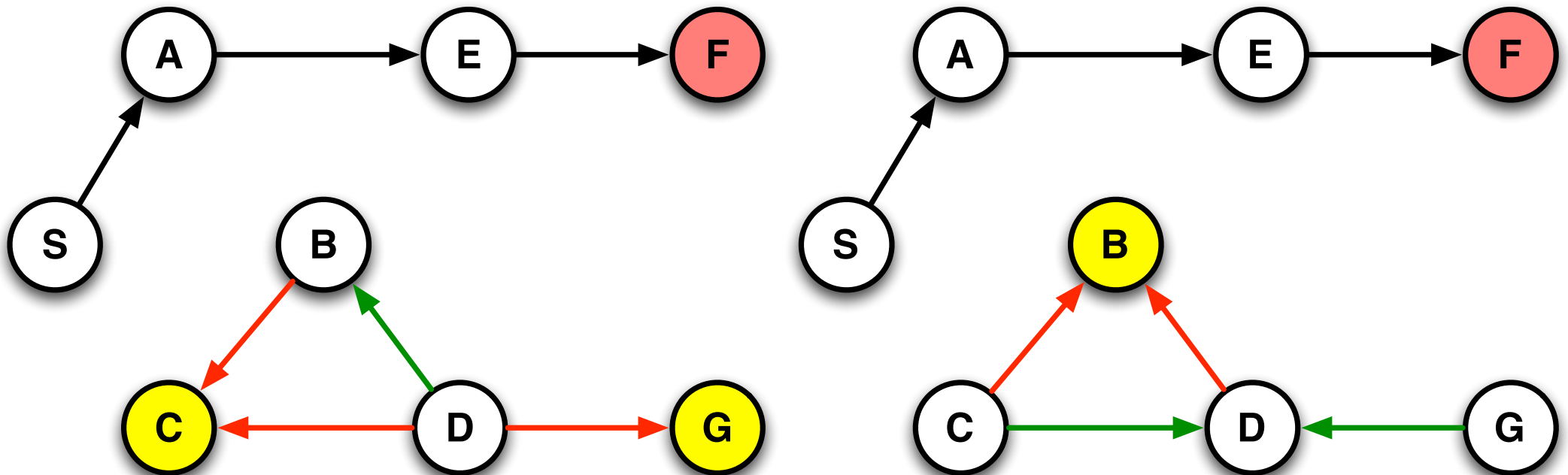
# Link Reversal if Network is Partitioned



# Link Reversal if Network is Partitioned



# Link Reversal if Network is Partitioned



# TORA

- **Separate network does not terminate with link reversals**
- **Mechanism for recognizing partitioning**
  - TORA (Temporally-Ordered Routing Algorithm (TORA))
    - Park, Corson, Highly Adaptive Distributed Routing Algorithm for Mobile Wireless Networks, Infocom 1997
  - Analysis of link reversal provides this information

# Link Reversal

- **Reactive protocol**
  - Repair only when data packet is not delivered
- **Proactive protocol**
  - Hello packets check all connections
- **Link reversal can be both proactive and reactive**

# Routing in MANETs

## ► Routing

- Determination of message paths
- Transport of data

## ► Protocol types

- proactive
  - Routing tables with updates
- reactive
  - repair of message paths only when necessary
- hybrid
  - combination of proactive and reactive

# Routing Protocols

## ► Proactive

- Routes are **demand independent**
- Standard Link-State und Distance-Vector Protocols
  - Destination Sequenced Distance Vector (**DSDV**)
  - Optimized Link State Routing (**OLSR**)

## ► Reactive

- Route are determined when needed
  - Dynamic Source Routing (**DSR**)
  - Ad hoc On-demand Distance Vector (**AODV**)
  - Dynamic MANET On-demand Routing Protocol
  - Temporally Ordered Routing Algorithm (**TORA**)

## ► Hybrid

- combination of reactive und proactive
  - Zone Routing Protocol (**ZRP**)
  - Greedy Perimeter Stateless Routing (**GPSR**)

# Optimized Link State Routing

## ► Literature

- RFC3626: Clausen, Jacquet, *Optimized Link State Routing Protocol*, 2003
- First published 1999

## ► Most proactive protocols are based on

- Link-state routing
- Distance-Vector routing



# Link State Routing

- **Connections are periodically published throughout the network**
- **Nodes propagate information to their neighbors**
  - i.e. flooding
- **All network information is stored**
  - with time stamp
- **Each node computes shortest paths**
  - possibly also other route optimizations

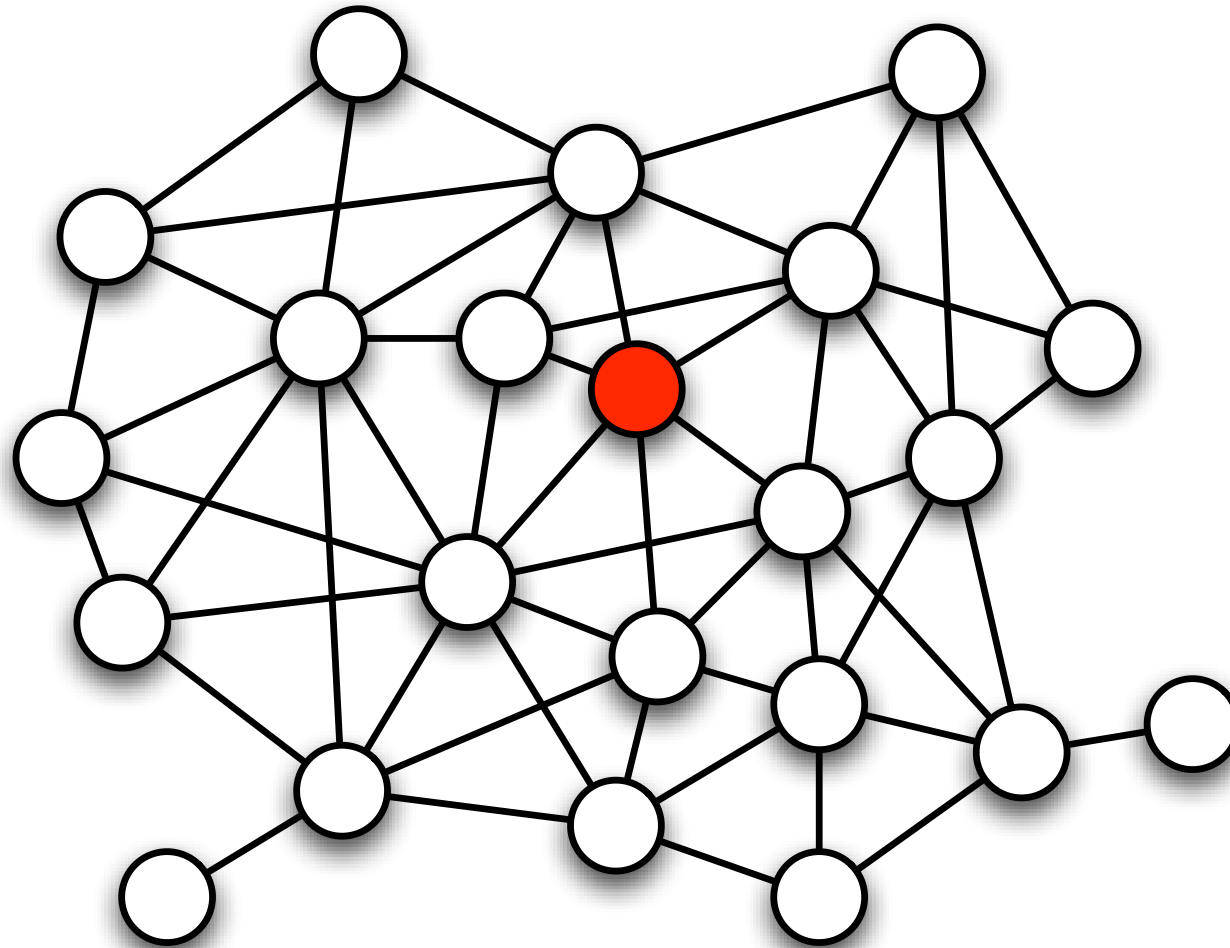
# Optimized Link State Routing (OLSR)

- **Each nodes broadcasts its neighborhood list**
  - Each node can determinate its 2-hop neighborhood
- **Reducing the number of messages**
  - fewer nodes participate in flooding
- **Multipoint relay node (MPRs)**
  - are chosen such that each node has at least one multipoint relay node as in its 2-hop neighborhood
  - Only multipoint relay nodes propagate link information
- **Node sends their neighborhood lists**
  - such that multipoint relay nodes in the 2-hop neighborhood can be chosen

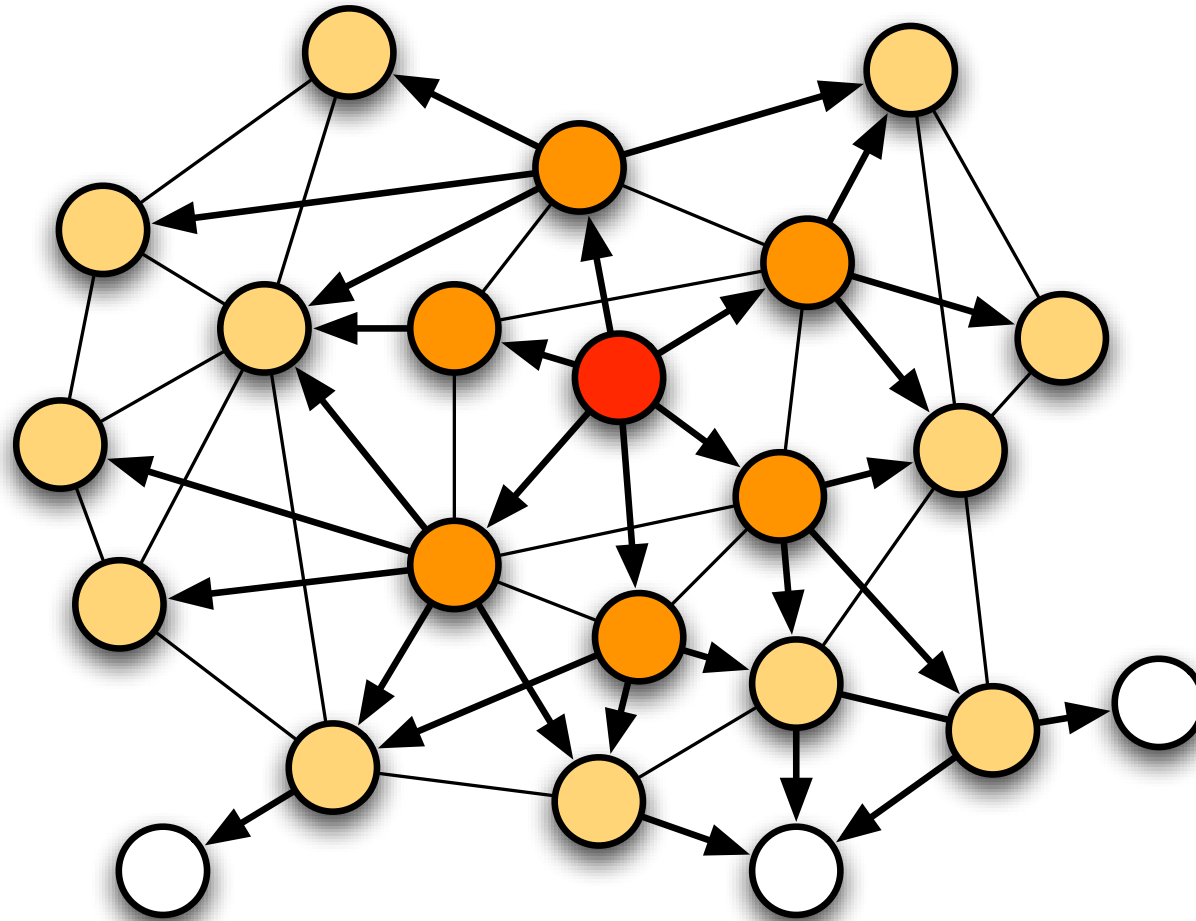
# Optimized Link State Routing (OLSR)

- **Combines Link-State protocol and topology control**
- **Topology control**
  - Each node chooses a minimal dominating set of the 2 hop neighborhood
    - ***multipoint relays (MPR)***
    - Only these nodes propagate link information
    - More efficient flooding
- **Link State component**
  - Standard link state algorithm on a reduced network

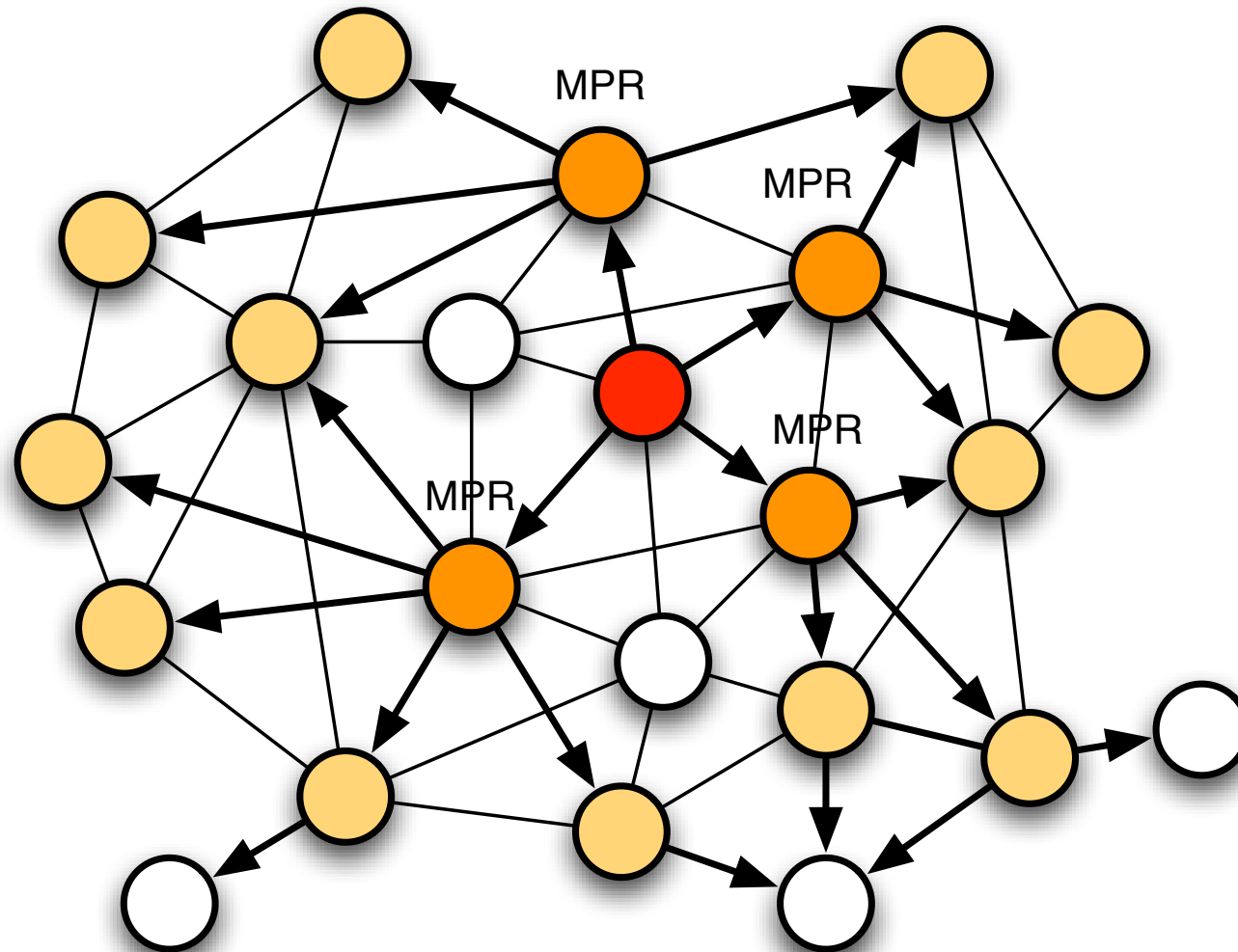
# Optimized Link State Routing (OLSR)



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# Selection of MPRs

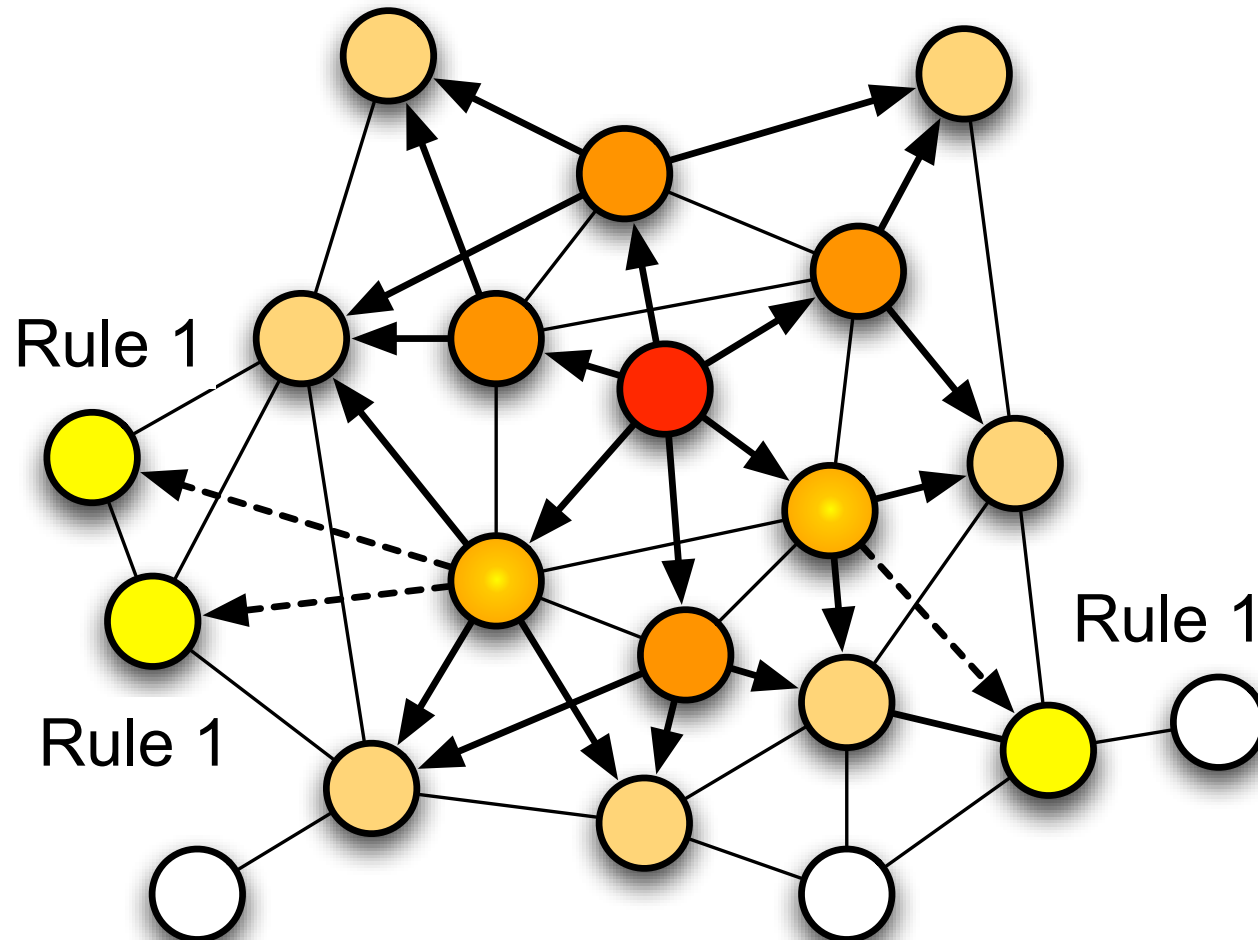
- **Multipoint Relaying for Flooding Broadcast Messages in Mobile Wireless Networks, Amir Qayyum, Laurent Viennot, Anis Laouiti, HICCS 2002**
- **Problem is NP-complete**
- **Heuristics**
  - recommended for OLSR
- **Notations**
  - $N(x)$ : 1 hop neighborhood of  $x$
  - $N^2(x)$ : 2 hop neighborhood of  $x$
  - Alle connections are symmetrical

# Selection of MPRs

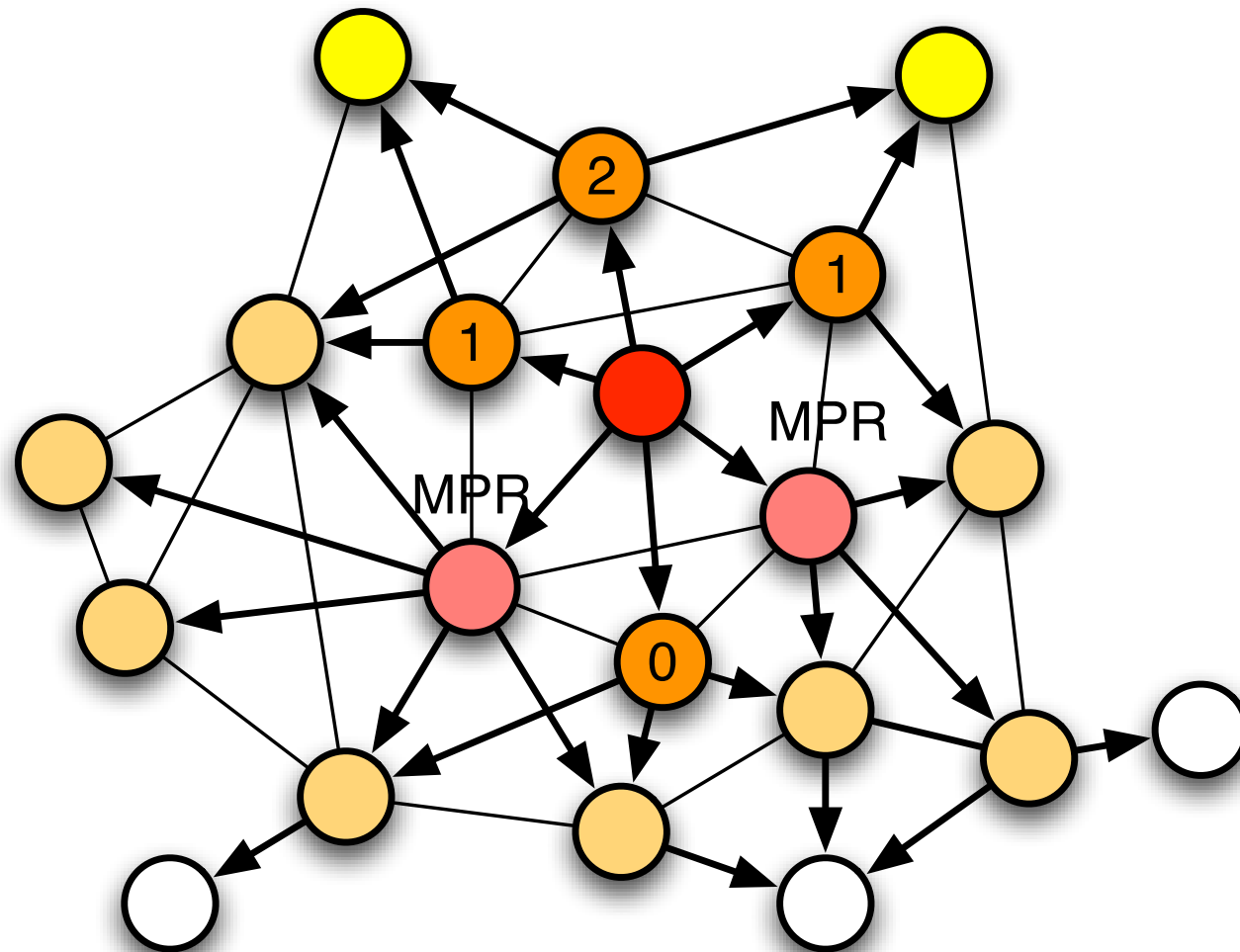
- ▶ **At the beginning there is no MPR**
  - Each node chooses its MPRs
- ▶ **Rule 1: A node of  $x$  is selected as MPR, if**
  - it in  $N(x)$  and
  - it is the only neighborhood node in the node  $N^2(x)$
- ▶ **Rule 2: If nodes in  $N^2(x)$  are not covered:**
  - Compute for each node in  $N(x)$  the number of uncovered nodes in  $N^2(x)$
  - Select as MPR the node that maximizes the value



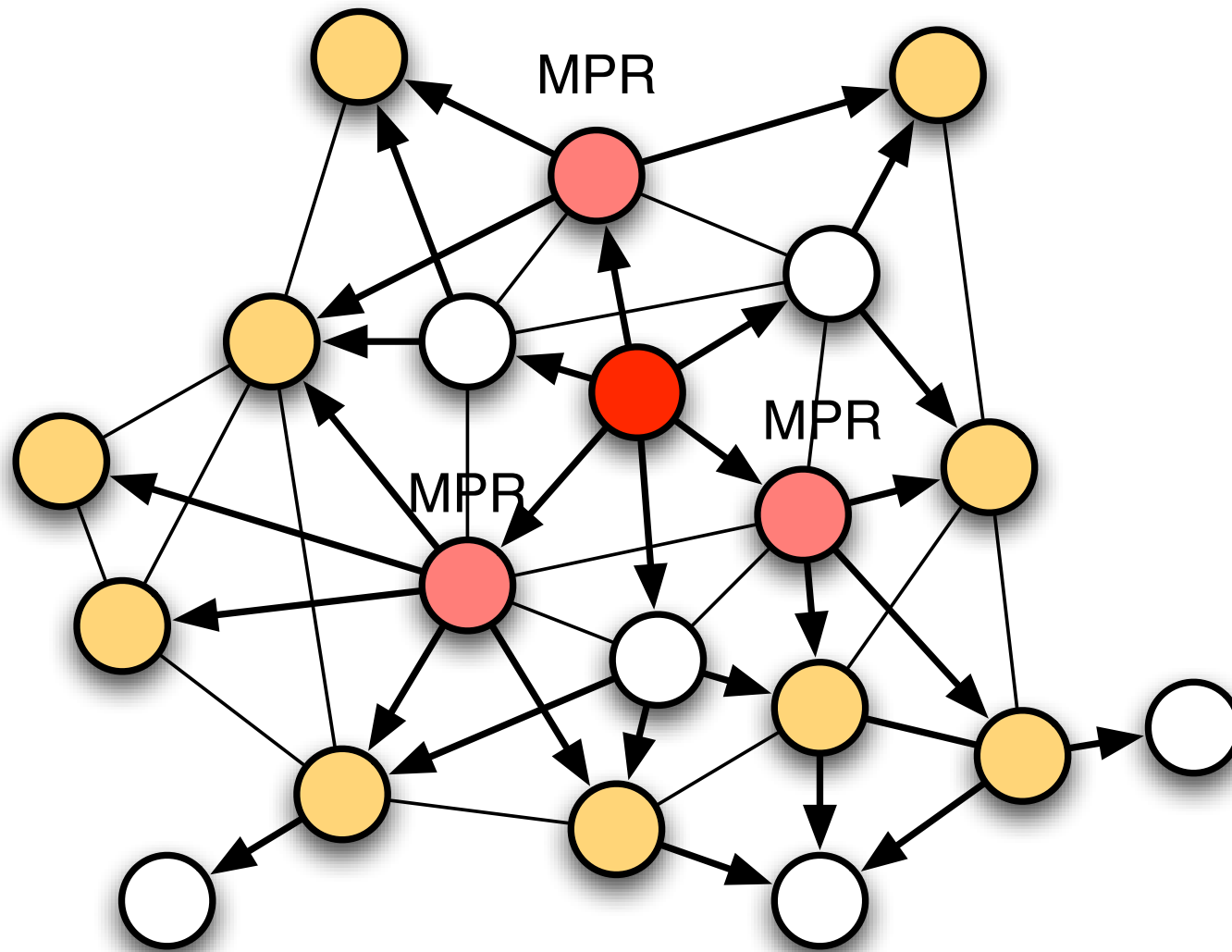
# Rule 1



# Rule 2



# MPRs



# OLSR

- **OLSR is flooding link information using MPRs**
  - Multipoint-Relays
- **Receivers choose their own MPRs for propagating**
  - Each node chooses its own MPRs
- **Routes use only MPRs as intermediate nodes**

# Zone Routing Protocol (ZRP)

## ‣ Haas 1997

- *A new routing protocol for the reconfigurable wireless networks*, Proc. of IEEE 6th International Conference on Universal Personal Communications, 562–566

## ‣ Zone Routing Protocol combine

- Proactive protocol
  - for local routing
- reactive protocol
  - for global routing

# ZRP

- **Routing zone of a node  $x$** 
  - Nodes in a given maximum hop-distance  $d$
- **Peripheral nodes**
  - all nodes have exactly the hop-distance  $d$  have
  - within the routing zone  $x$

# ZRP

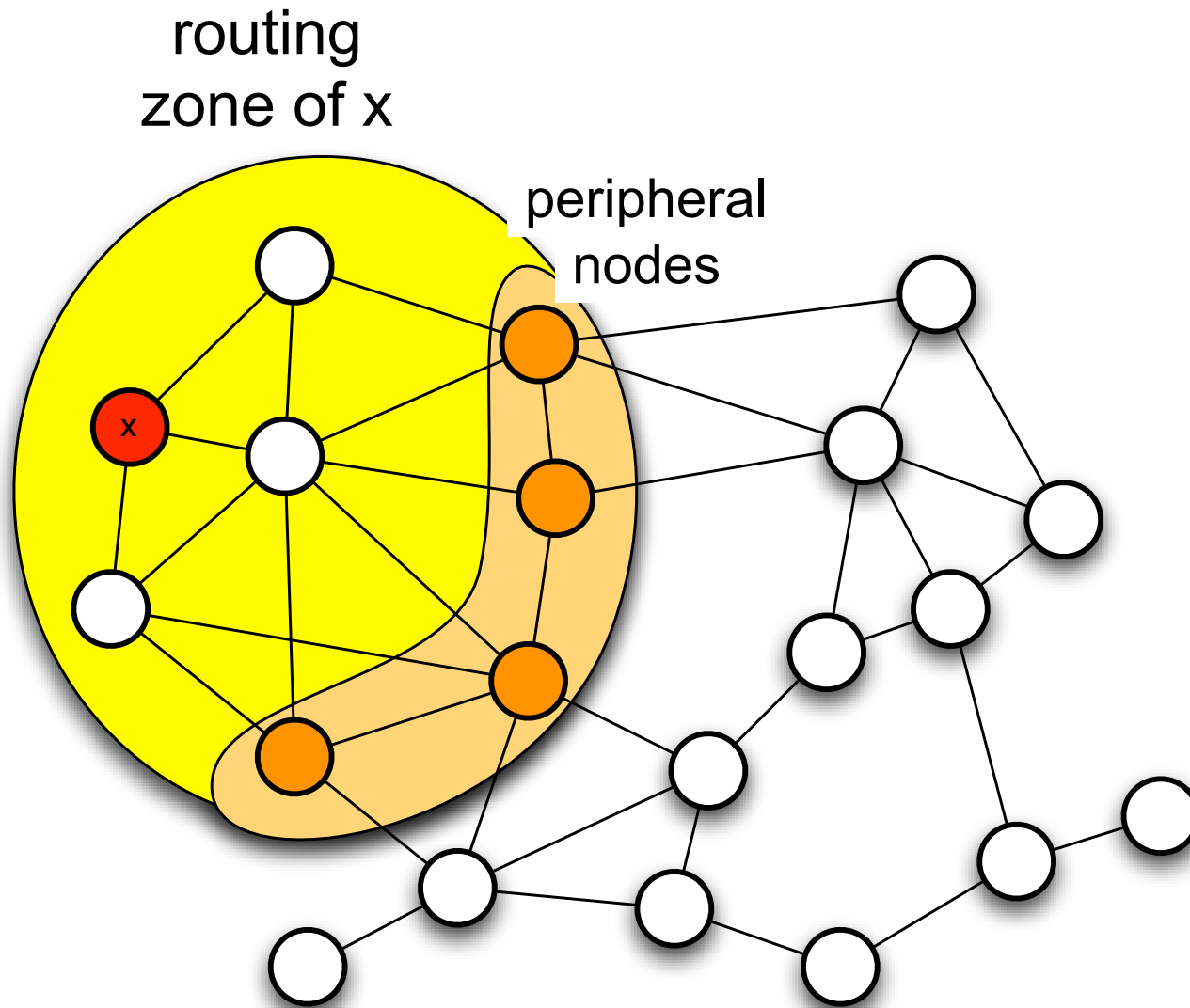
## ► Intra zone routing

- proactive update the connection information in the routing zone of node
  - e.g. with link state or distance vector protocols

## ► Inter zone routing

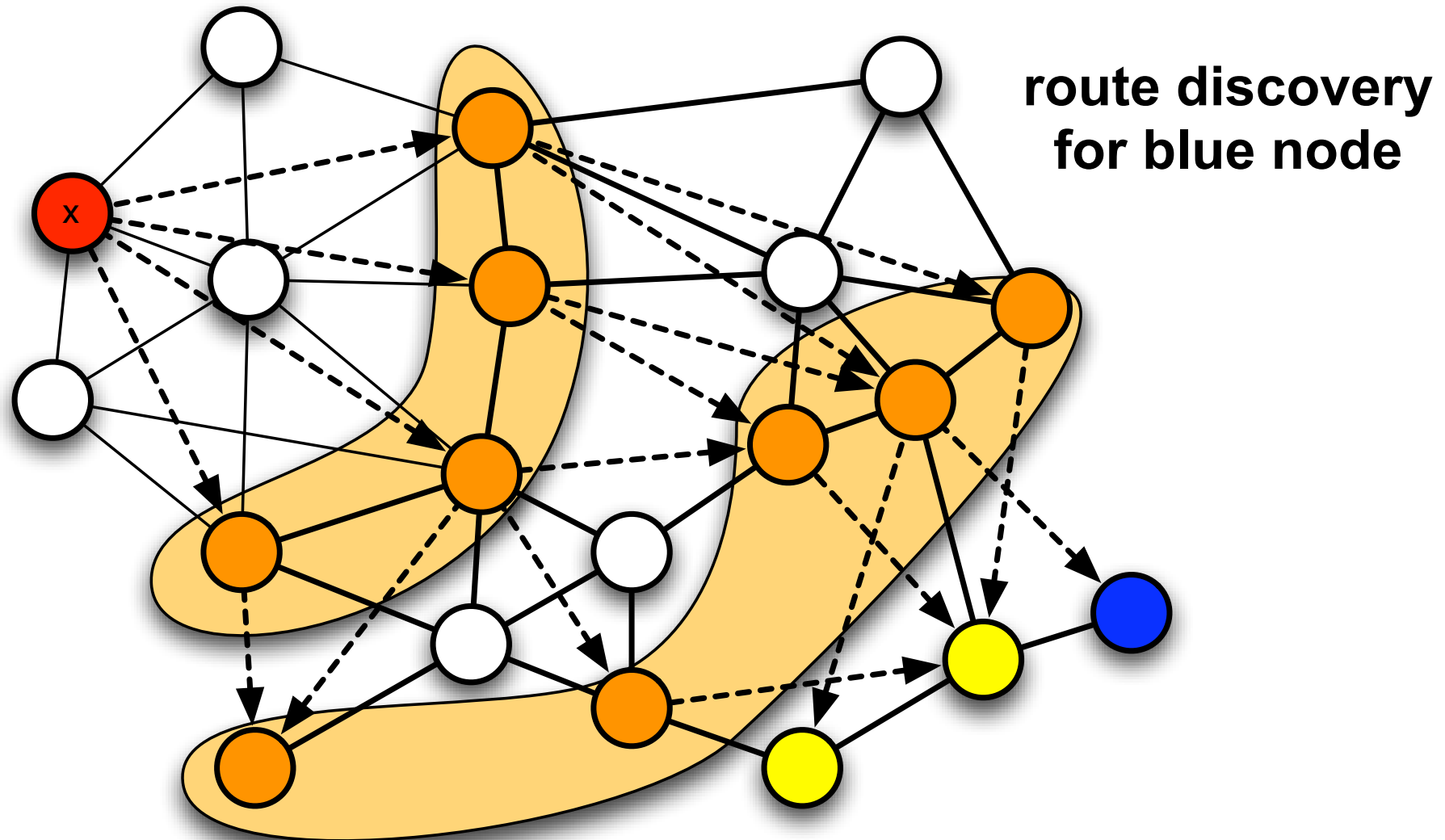
- Reactive route discovery is used for distant / unknown nodes
- Procedure similar to DSR
- Only peripheral nodes reach further information

# ZRP: Example with radius $d=2$

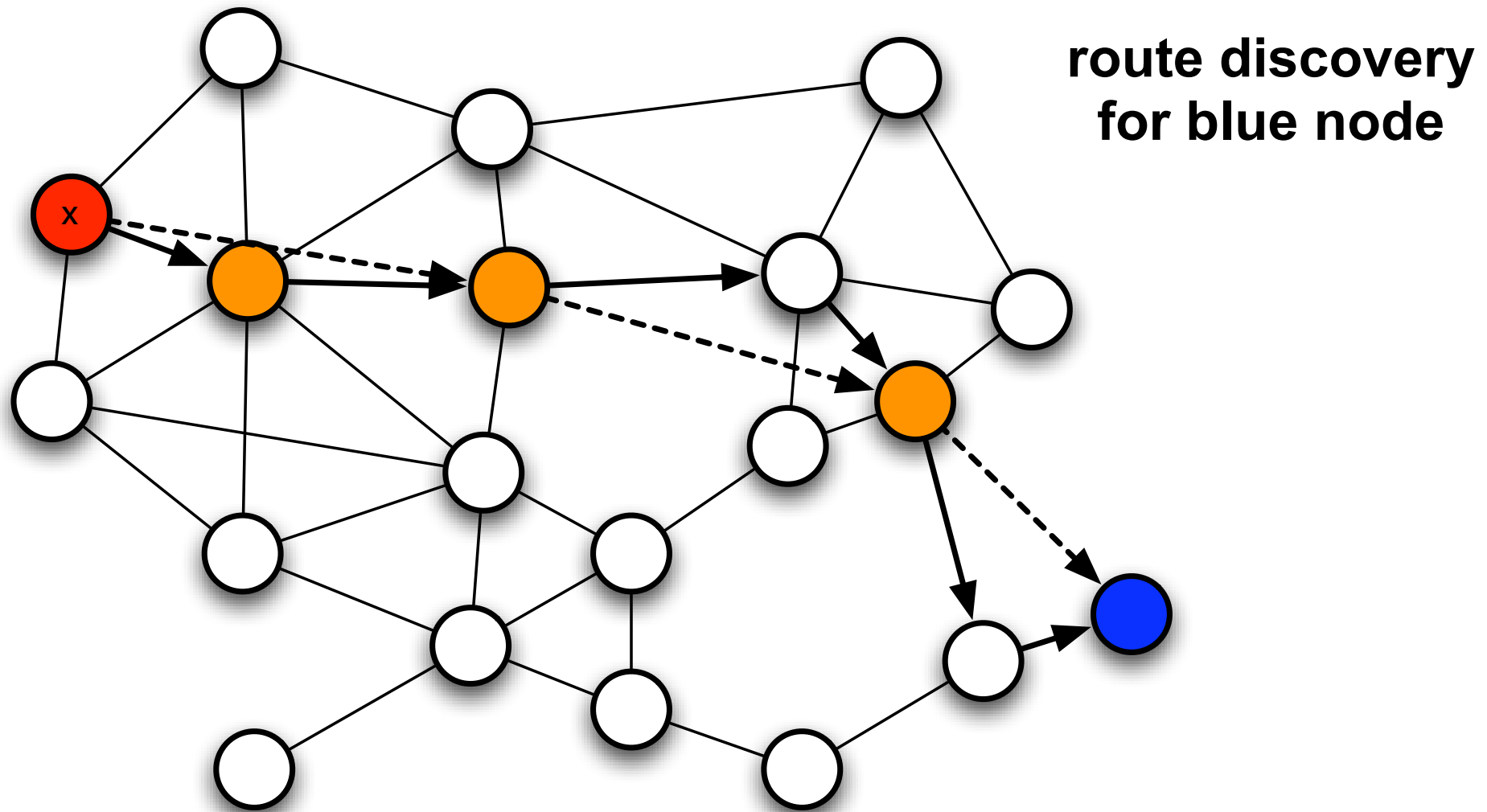




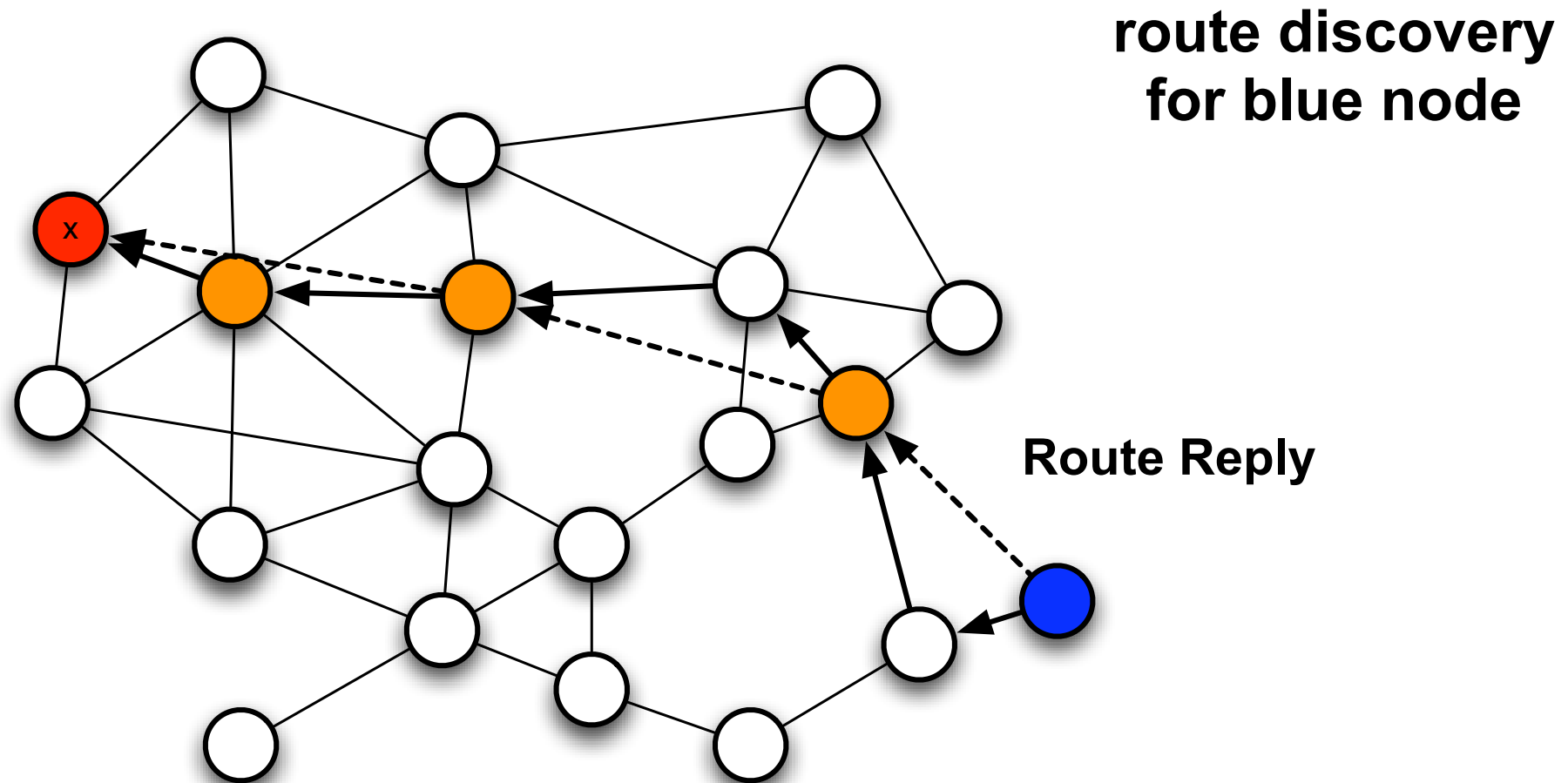
# ZRP: Example with radius $d=2$



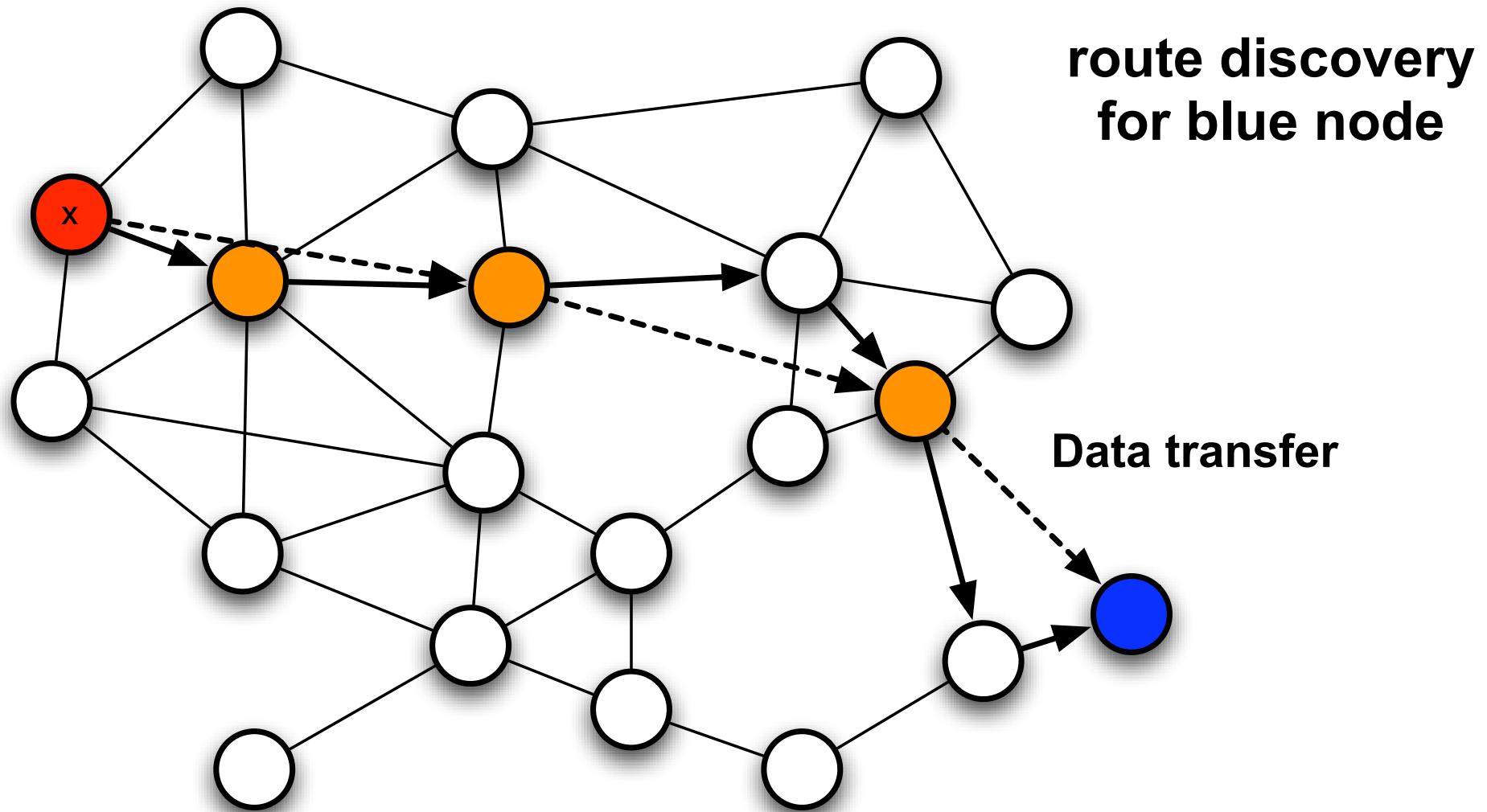
# ZRP: Example with radius $d=2$



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# ZRP: Example with radius $d=2$





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