

Algorithms for Radio Networks

MACAW

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MACAW

- Bharghavan, Demers, Shenker, Zhang
 - MACAW: A Media Access Protocol for Wireless LAN's, SIGCOMM 1994
 - Palo Alto Research Center, Xerox

→ Aim

- Redesign of MACA
- Improved backoff
- Fairer bandwidth sharing using Streams
- Higher efficiency
 - by 4- and 5-Handshake

Acknowledgment in the Data Link Layer

MACA

- does not use Acks
- initiated by Transport Layer
- very inefficient
- How can MACA use Acks?

MACAW 4 Handshake

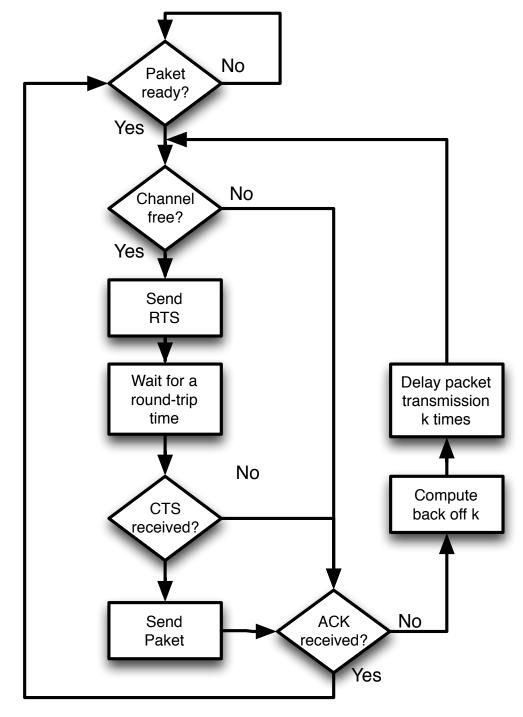
Participants

- Sender sends RTS
- Receiver answers with CTS
- Sender sends data packet
- Receiver acknowledges (ACK)

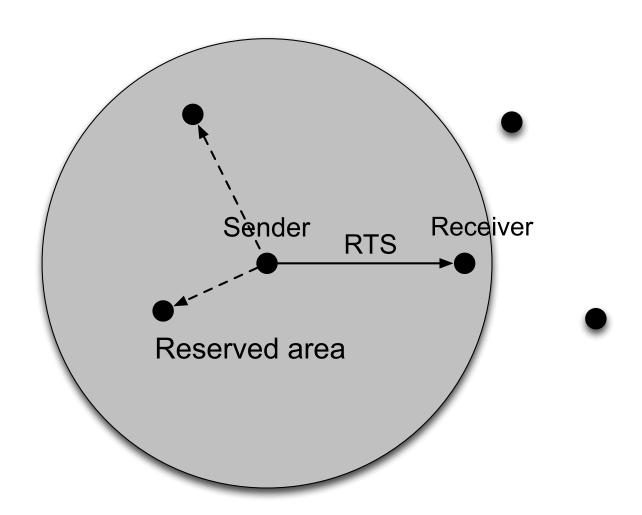
Third parties

- Nodes receiving RTS or CTS are blocked for some time
- RTS and CTS describe the transmission duration
- Sender repeats RTS, if no ACK has been received
 - If receiver has sent ACK
 - then the receiver sends (instead of CTS) another ACK

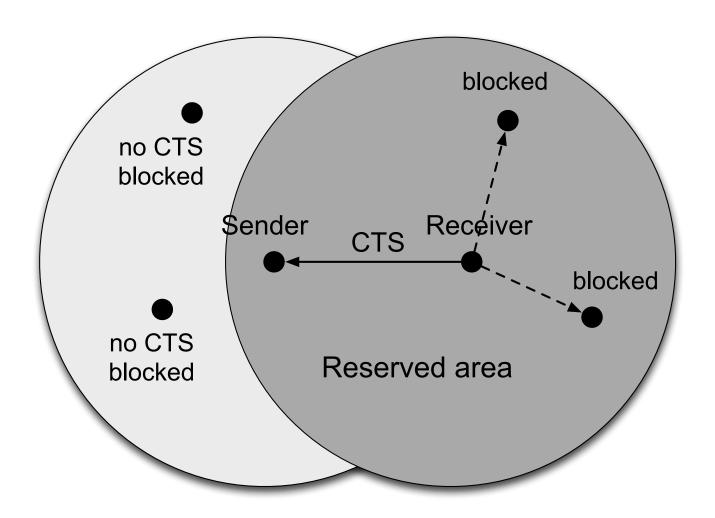
MACAW



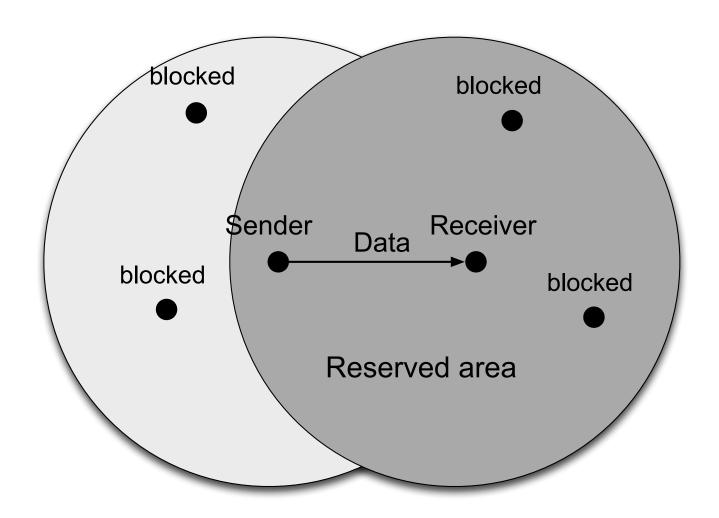
MACA 4-Handshake RTS



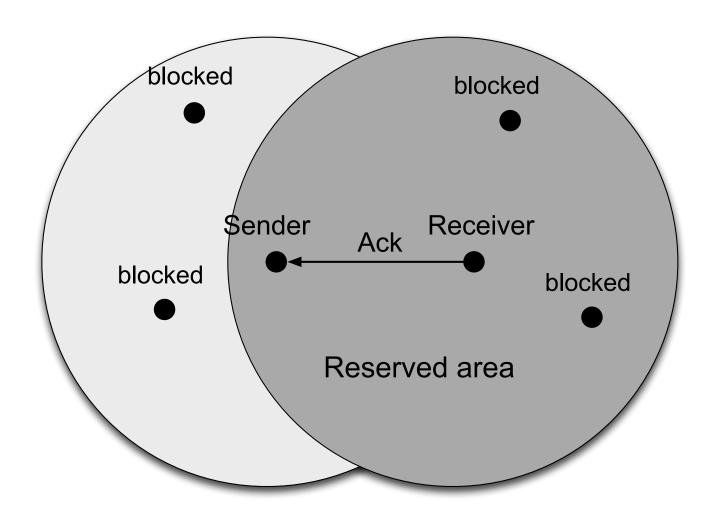
MACAW 4-Handshake CTS



MACAW 4-Handshake Data



MACAW 4-Handshake Ack



Acknowledgments

- Adding ACKs to MACA
 - In MACA done by transport layer
- leads to drastical improvements of throughput even for moderate error rates

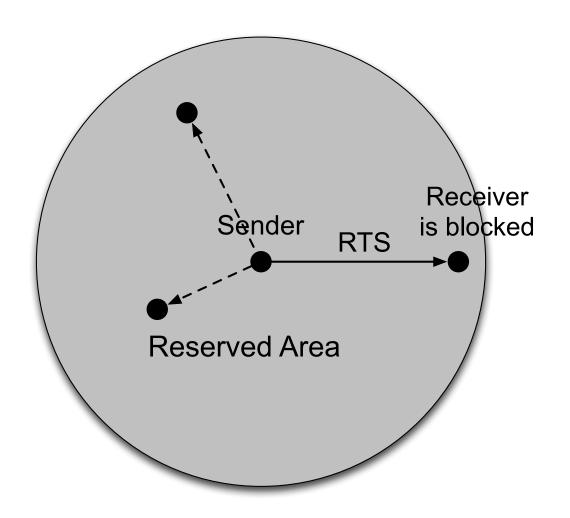
	throughput	
error rate	RTS-CTS- DATA	RTS-CTS- DATA-ACK
0	40	37
0,001	37	37
0,01	17	36
0,1	2	10

MACAW 4 Handshake

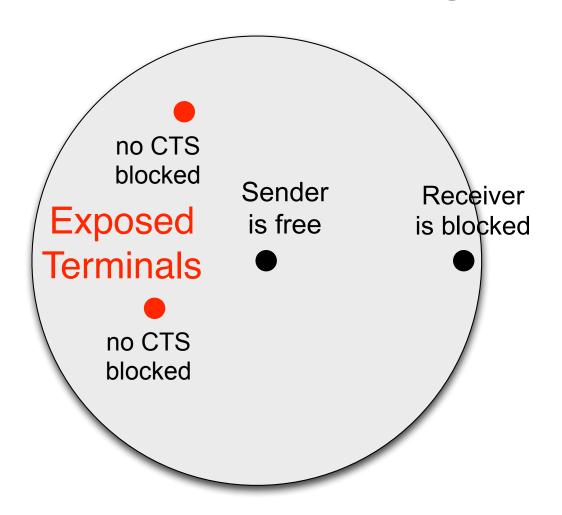
Worst-Case blockade

- Sender sends RTS
- Receiver is blocked
- Sender is free
- But the environment of the sender is blocked

MACAW 4-Handshake RTS



MACAW 4-Handshake CTS is missing



MACAW 5 Handshake

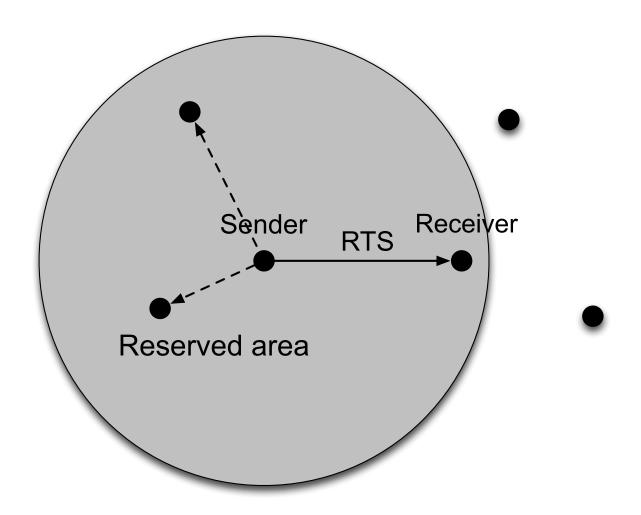
- 4-Handshake increases Exposed Terminal Problem
 - Overheard RTS blocks nodes
 - even if there is no data transfer
- Solution
 - Exposed Terminals are informed whether data transmission occurs
 - Short message DS (data send)
- 5 Handshake reduces waiting time for exposed terminals

MACAW 5 Handshake

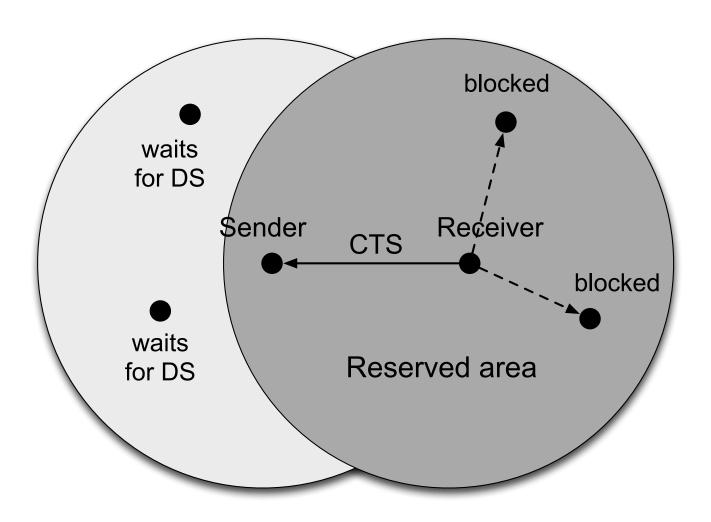
Participants

- Sender sends RTS
- Receivers answers with CTS
- Sender sends DS (Data Send)
- Sender sends DATA PACKET
- Receiver acknowledges (ACK)
- RTS and CTS announce the transmission duration
- Blocked nodes
 - have received RTS and DS
 - have received CTS
- Small effort decreases the number of exposed terminals

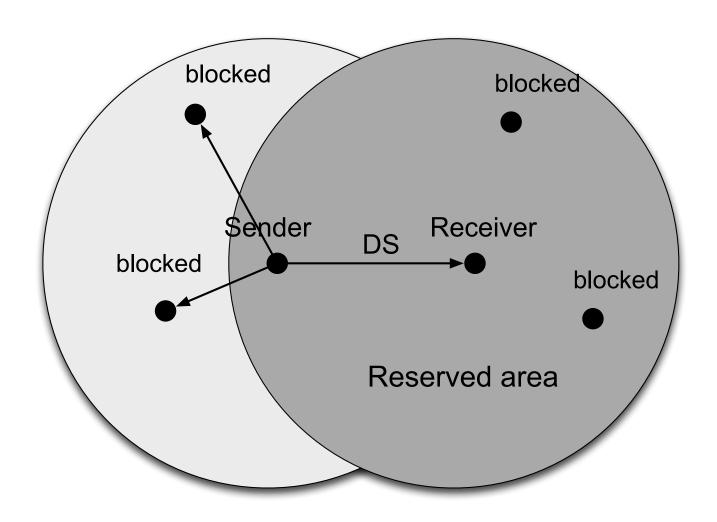
MACAW 5-Handshake RTS



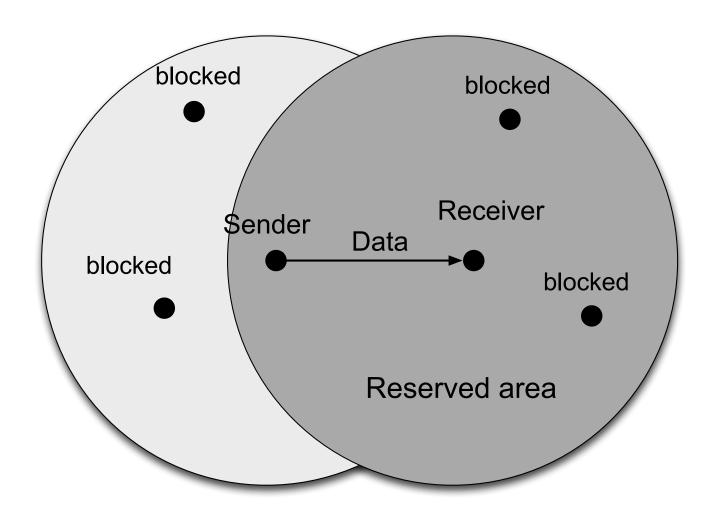
MACAW 5-Handshake CTS



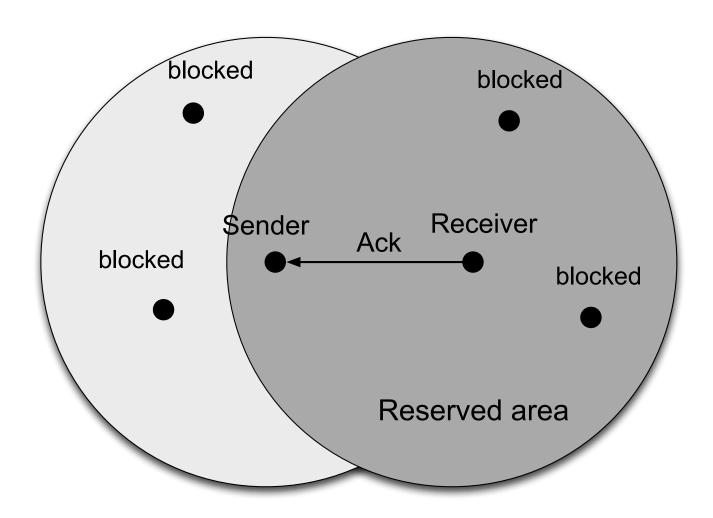
MACAW 5-Handshake DS



MACAW 5-Handshake Data



MACAW 5-Handshake ACK



Unfair Distribution

- 4 and 5-Handshake create unfair distribution
 - A has a lot of data for B
 - D has a lot of data for C
 - C receives B and D, but does not receive A
 - B can receive A and C, but does not hears D

- A is the first to get the channel
- D sends RTS and is blocked
 - Backoff of D is doubling
- At the next transmission
 - A has smaller backoff
 - A has higher chance for next channel access



RRTS

Solution

- C sends RRTS (Request for Request to Send)
 - if ACK has been received
- D sends RTS, etc.

Why RRTS instead of CTS?

- If neighbors receive CTS, then they are blocked for a long time
- Possibly, D is not available at the moment



Backoff Algorithms

- After collision wait random time from {1,.. Backoff}
- Binary Exponential Backoff (BEB) algorithm
 - Increase after collision
 - backoff = min{2 backoff, maximal backoff}
 - Else:
 - backoff = Minimal Backoff
- Multiplicative increase, linear decrease (MILD)
 - Increase:
 - backoff = min{1.5 backoff, maximal backoff}
 - Else:
 - backoff = max{backoff 1, minimal-backoff}

Information Dissemination for Backoff-Algorithm

Backoff parameter are overheard

- participants adapt the parameters to the overheard backoff values
- using MILD

Motivation

 if a participant has the same backoff value, then the fairness has been reached



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