5.1: Introduction

One of the most important prevalent problems in distributed systems

The Internet

- no central entity
- > 400,000 routers of AS (autonomous systems) use BGP (border gateway protocol) to determine routes between AS
- within AS distributed routing protocols update routing tables

Mobile Ad-Hoc Networks

- no central control
- dynamic wireless nodes
- connections appear and disappear at any time
5.2: Routing in the Internet

- Packet forwarding and route selection
- Each node has a routing table
- When a packet needs to be processed, choose the best choice from the table
- Decrease TTL (time to live counter)
- If TTL = 0 then delete packet

- Route selection
  - Programming of the routing table
  - Originally: static (manually)
  - Always single-path
  - Flat versus hierarchical
  - Intradomain and interdomain
5.3: Routing Paradigms

Optimization goals
- delay
- hop count
- throughput
- reliability
- monetary cost

Routing decisions
- Distributed routing
  - routers decide
  - BGP, Distance-Vector, Link-State Routing
- Source routing
  - senders decide
  - DSR (Dynamic Source Routing), Onion Routing (TOR)
- Centralized routing
  - one (possibly external) instance decides all routing
  - (mobile) phone network, static routing
Adaptivity

- Deterministic routing algorithm
  - optimizes the path (usually the hop count)
  - each packet is using the same route

- Oblivious routing algorithms
  - oblivious to the status of the network
  - Valiant has proved that in a two-dimensional network randomized oblivious routing optimizes the throughput
  - packets are routed in random or cyclic directions

- Adaptive routing algorithms
  - use information about network traffic / channel status
  - avoid congested areas
  - avoid unreliable links
Switching Techniques

- **Circuit Switching**
  - First, determine a path and reserve the router nodes
  - Then, send the packets on this path exclusively reserved for this connection

- **Packet Switching**
  - store-and-forward switchings
  - each packet is *individually* routed from source to target

- **Virtual cut-through switching (VCT)**
  - Messages are split into packets
  - if the channel to the next router is free, the packet is immediately forwarded
  - otherwise the packets are buffered in the router with the first blockade and sent if the channel is free again

- **Wormhole (WH) switching**
  - like VCT but routers have small buffers
  - packet string is stored on buffers like a snake
5.4: Bellman Ford

The Bellman-Ford Algorithm computes the shortest path problem towards \( t \) from each node in graph \( G = (V, E) \) with weights \( w(u, v) \).

**Init-Target\((G, w, t)\)**
- **Init-Target\((G, w)\)**
  - for all \( v \in V \) do
    - \( d(v) \leftarrow \infty \)
    - \( \pi(v) \leftarrow v \)
  - \( d(t) \leftarrow 0 \)

**Relax\((u, v)\)**
- **Relax\((u, v)\)**
  - if \( d(u) > w(u, v) + d(v) \) then
    - \( d(u) \leftarrow w(u, v) + d(v) \)
    - \( \pi(u) \leftarrow v \)
5. Distributed Routing

5.4. Bellman Ford

Init-Target

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Relax
Relax
Relax
Init-Target for Directed Graphs

![Diagram of directed graphs with node labels and directed edges with weights. The graph on the left shows initial weights, and the graph on the right shows the updated weights after one iteration of the Bellman-Ford algorithm. The target node is highlighted in red.]

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Relax for Directed Graphs
Compute for each (target) node of the network the following.

**Distributed Bellman Ford for target $t$**

*(Distance-Vector Routing)*

- If node is $t$ then $d(t) \leftarrow 0$; $\pi(t) \leftarrow t$
- If a message from $u$ to $\pi(u)$ fails then $d(u) \leftarrow \infty$
- If $u$ detects a new neighbor $v$ then send $(u, d(u))$ to $v$
- If $u$ receives $(v, d(v))$ from $v$
  - if $d(u) > d(v) + w(u, v)$ or $v = \pi(u)$ then
    - $d(u) \leftarrow d(v) + w(u, v)$
    - $\pi(u) \leftarrow v$
  - if $d(u) = \infty$ then $\pi(u) \leftarrow u$
- Every time $d(u)$ or $\pi(u)$ has changed $u$ sends $(u, d(u))$ to all neighbors
Distance Vector
Correctness

- Let $|u, v|$ be the distance from $u$ to $v$
- Assume that the weights are constant
- Then, at each time of the operation $d(u) > |u, t|$
- If the shortest path from $u$ to $t$ is $(u, v_1, v_2, \ldots, v_n, t)$ and
  - a message is sent from $t$ to $v_n$ and then from $v_n$ to $v_{n-1}$, etc.
  - then $d(v_n) = |v_n, t|$, $d(v_{n-1}) = |v_{n-1}, t|$, \ldots, $d(u, t) = |u, t|$
- So, for each shortest path of finite length eventually Distributed Bellman Ford converges
Problems of Distributed Bellman Ford

- **Negative path cycles**
  - Bellman-Ford works fine as long as the shortest path is finite
  - if a negative path cycle exist, then the shortest path is infinite

- **Dynamic graphs**
  - Temporarily wrong routes
    - During the distributed computation messages might take wrong directions
    - could even revisit the same node more than once
  - Lost connections
    - Bad news travels slowly (or not at all)
    - If the distance increases then no messages are produced
    - Old an wrong information is preserved

- Temporarily undefined routes
  - For speeding up store last message of each neighbor
  - Use this information when new messages arrive

- **Count-to-Infinity Problem**
  - If the target is not reachable and at least two nodes remain connected
  - then the distance is updated towards infinity
Negative Path Cycles

![Diagram of Negative Path Cycles]

The diagram illustrates a network with nodes connected by weighted paths. The target node is highlighted, and the network contains negative cycle paths, which indicate paths where the sum of the weights is negative.
Temporarily Wrong Routes
Count to Infinity Problem

- If the target is not reachable and at least two nodes remain connected
- then the distance is updated towards infinity
Solutions to the Count-to-Infinity Problem

- **Split Horizon**
  - A node does not advertise routes to nodes on its path to the target

- **Poison reverse**
  - A variant of split horizon
  - A node does not advertise routes to nodes on its path to the target
  - With the exception of a value of $\infty$ is advertised immediately towards target

- **Link State Routing**
  - All link information is broadcast in the network
  - Each node computes the shortest path on this information

- **Path Vector Protocol**
  - Instead of only storing the next link
  - Store the complete path and forward it to the neighbors
  - Used in BGP
5. Distributed Routing

5.4. Bellman Ford

Split Horizon

Poison Reverse

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Distance-Vector Routing with Split Horizon for target $t$

- If node is $t$ then $d(t) \leftarrow 0$; $\pi(t) \leftarrow t$
- If a message from $u$ to $\pi(u)$ fails then
  - $d(u) \leftarrow \infty$
  - $\pi(u) \leftarrow u$
- If $u$ detects a new neighbor $v$ then
  - send $(u, d(u))$ to $v$
- If $u$ receives $(v, d(v))$ from $v$
  - if $d(u) > d(v) + w(u, v)$ or $v = \pi(u)$ then
    - $d(u) \leftarrow d(v) + w(u, v)$
    - $\pi(u) \leftarrow v$
- Periodically and every time $d(u)$ or $\pi(u)$ has changed $u$ sends $(u, d(u))$ to all neighbors except for $\pi(u)$

Split horizon rule: information is not sent towards target
Distance-Vector Routing with Poison Reverse for target $t$

- If node is $t$ then $d(t) \leftarrow 0$; $\pi(t) \leftarrow t$
- If a message from $u$ to $\pi(u)$ fails then
  - $d(u) \leftarrow \infty$
  - $\pi(u) \leftarrow u$
- If $u$ detects a new neighbor $v$ then
  - send $(u, d(u))$ to $v$
- If $u$ receives $(v, d(v))$ from $v$
  - if $d(u) > d(v) + w(u, v)$ or $v = \pi(u)$ then
    - $d(u) \leftarrow d(v) + w(u, v)$
    - $\pi(u) \leftarrow v$
- Periodically and every time $d(u)$ or $\pi(u)$ has changed send $(u, d(u))$ to all neighbors except for $\pi(u)$
- If $d(u)$ has changed to $\infty$ then send $(u, d(u))$ to $\pi(u)$

Poison reverse: remove loops before they can propagate
Path Vector Protocol for target $t$

- $d(t) \leftarrow 0$; $p(t) \leftarrow (t)$
- $t$ send $s(t, d(t), p(t))$ to all neighbors once
- If a message from $u$ to the first node of $p(u)$ fails then
  - $d(u) \leftarrow \infty$
  - $p(u) \leftarrow ()$
- If $u$ detects a new neighbor $v$ then
  - send $(u, d(u), p(u))$ to $v$
- If $u$ receives $(v, d(v), p(v))$ from $v$
  - if $d(v) = \infty$ and $v \in p(u)$ then $u$’s route has vanished
    - $d(u) \leftarrow \infty$
    - $p(u) \leftarrow ()$
  - if $u \not\in p(v)$ and $d(u) > d(v) + w(u, v)$ then only if $u$ is not in the
    - $d(u) \leftarrow d(v) + w(u, v)$
    - $p(u) \leftarrow (u, p(v))$
- Periodically and every time $d(u)$ or $p(u)$ has changed send $(u, d(u), p(u))$ to all neighbors
Relax in Path Distance Vector
Lost Target in Path Distance Vector

\begin{align*}
\text{target} & \quad (v, 4, p(v)) \\
v & \quad \infty \quad u \quad 2 \\
(\infty, \infty, ()) & \quad (v, \infty, ()) \\
4 & \quad v \quad 3 \\
0 & \quad \infty \quad v \quad 2 \\
\end{align*}
Applications of Routing Algorithms

- Mobile Ad Hoc Networks
  - reactive versus proactive protocols
  - proactive protocols like DV construct routing tables even without packets
  - when the network is too dynamic this is an overhead
  - then reactive protocols like Flooding, DSR, AODV are more adequate

- DV routing is used in Intra-Domain Routing
  - RIP (Routing Information Protocol)
  - EIGRP (Enhanced Interior Gateway Routing Protocol)

- main alternative for Intra-Domain Routing is Link-State-Routing used in
  - OSPF (Open Shortest Path First)
  - IS-IS (Intermediate System to Intermediate System Protocol)

- Path-Vector-Protocol is the worldwide standard in Inter-Domain Routing
  - BGP (Border Gateway Protocol)
End of Section 5
\[(ACD)^* : -\infty\]

\[A \rightarrow C \rightarrow D : -10\]

\[A \rightarrow C \rightarrow D : -15\]

\[10 \rightarrow 20 \rightarrow 5 \rightarrow 10 \rightarrow -20\]
5. Distributed Routing

5.4. Bellman Ford

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